



Chapter 10: Nonregular Languages

We show several examples of nonregular languages: those that cannot be defined by regular expressions.

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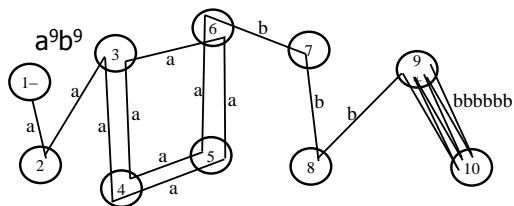


Example:

- $L = \{\Lambda, ab, aabb, aaabbb, aaaabbbb, \dots\}$
- $L = \{a^n b^n \mid n=0,1,2,3,4,\dots\}$
- $L = \{a^n b^n\}$
- $L \subset \text{language}(a^*b^*)$

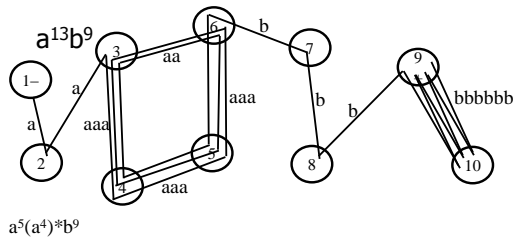
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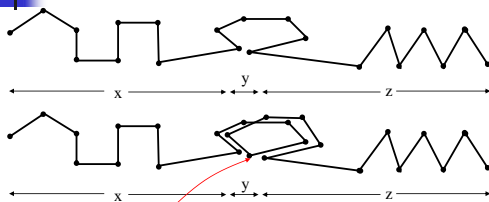


**The pumping lemma:**

Let L be any regular language that has infinitely many words. Then there exist three words x, y, z (where y is not the empty word) such that all words of the form:

$$xy^n z \quad n=1,2,3,4,\dots$$

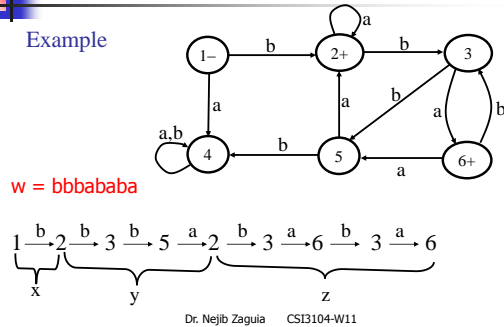
are in L .



first state that is visited more than once

$xz \in L, xyz \in L, xyxz \in L, xyxyz \in L, \dots, xy^nbz \in L$.

Example



Theorem: $L = \{a^n b^n \mid n=0,1,2,3,4,\dots\}$ is not regular.

EQUAL = all words with the same number of a's and b's.

EQUAL = $\{\Lambda, ab, ba, aabb, abab, abba, baab, baba, bbaa, \dots\}$.

Theorem: EQUAL is not regular.

Theorem: $L = \{a^n b a^n \mid n=0,1,2,3,4,\dots\}$
 $= \{b, aba, aabaa, \dots\}$ is not regular.

Pumping Lemma: (version 2)

Let L be any regular language that has infinitely many words that is accepted by a finite automaton with N states. All words w in L that have more than N letters can be decomposed into words x,y,z such that:

1. y is not the empty word
2. $\text{length}(x) + \text{length}(y) \leq N$
3. $w = xyz$
4. for all $n \geq 1$, $xy^n z$ is in L .



Theorem: PALINDROME is not regular.

Theorem: PRIME = $\{a^p \mid p \text{ is a prime number}\}$
= $\{aa, aaa, aaaaa, aaaaaa, \dots\}$
is not regular.
