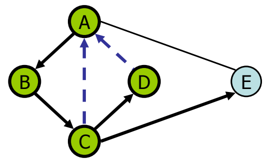
Graph Traversals

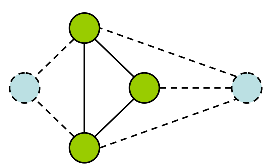


Outline and Reading

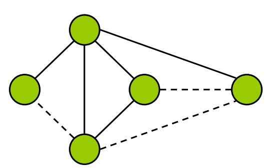
- Definitions (§12.1)
  - Subgraph
  - Connectivity
  - Spanning trees and forests

Subgraphs

- A subgraph S of a graph G is a graph such that
  - The vertices of S are a subset of the vertices of G
  - The edges of S are a subset of the edges of G
- A spanning subgraph of G is a subgraph that contains all the vertices of G



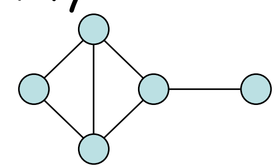
Subgraph



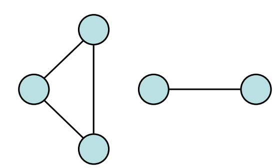
Spanning subgraph

Connectivity

- A graph is connected if there is a path between every pair of vertices
- A connected component of a graph G is a maximal connected subgraph of G



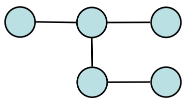
Connected graph



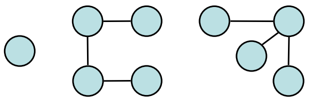
Non connected graph with two connected components

Trees and Forests

- A tree is an undirected graph T such that
  - T is connected
  - T has no cycles
- A forest is an undirected graph without cycles (a collection of trees).
- The connected components of a forest are trees



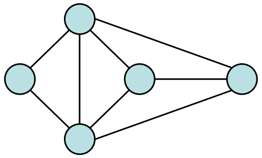
Tree



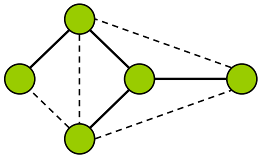
Forest

Spanning Trees and Forests

- A spanning tree of a connected graph is a spanning subgraph that is a tree
- A spanning tree is not unique unless the graph is a tree
- Spanning trees have applications to the design of communication networks
- A spanning forest of a graph is a spanning subgraph that is a forest



Graph



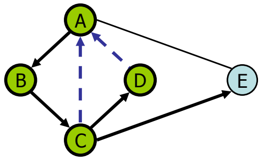
Spanning tree

Graph Traversals

A traversal of a graph *G*:

- Visits all the vertices and edges of *G*
- Determines whether *G* is connected
- Computes the connected components of *G*
- Computes a spanning forest of *G*
- Build a spanning tree in a connected graph

Depth-First Search



Outline and Reading

- Depth-first search (§12.3.1)
  - Algorithm
  - Example
  - Properties
  - Analysis
- Applications of DFS (§12.3.1)
  - Path finding
  - Cycle finding

Depth-First Search

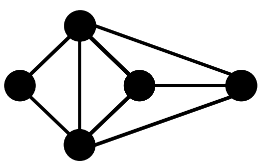
Depth-First Search is a graph traversal technique that:

- on a graph with n vertices and m edges takes $O(n + m)$ time
- can be further extended to solve other graph problems
  - Find and report a path between two given vertices
  - Find a cycle in the graph

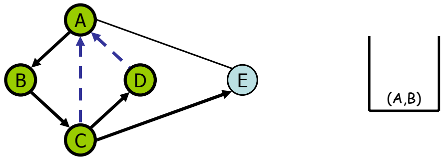
The idea:

Starting at an arbitrary vertex, follow along a simple path until you have get to a vertex which has no unvisited adjacent vertices.

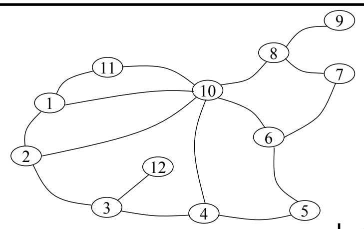
Then start tracing back up the path, one vertex at a time, to find a vertex with unvisited adjacent vertices.



DFS Algorithm – With a Stack



Example



Visited: {1} to visit

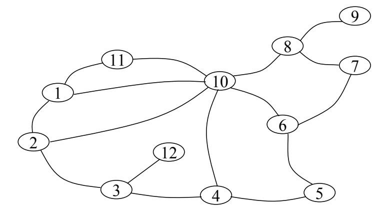
(1,2)
(1,11)
(1,10)

POP → (1,2)
Visited: {1,2} to visit
T={(1,2)}

(2,3)
(2,10)
(1,11)
(1,10)

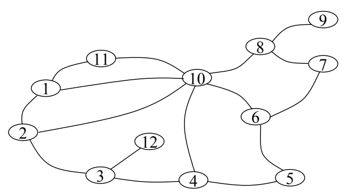
POP → (2,3)
Visited: {1,2,3} to visit
T={(1,2), (2,3)}

(3,12)
(3,4)
(2,10)
(1,11)
(1,10)



(3,4)
(2,10)
(1,11)
(1,10)

POP → (3,12)
Visited: {1,2,3,12} to visit
T={(1,2), (2,3), (3,12)}



(4,10)
(4,5)
(2,10)
(1,11)
(1,10)

POP → (3,4) to visit
Visited: {1,2,3,12,4}
T={(1,2), (2,3), (3,12), (3,4)}

POP → (4,10) to visit
Visited: {1,2,3,12,4,10}
T={(1,2), (2,3), (3,12), (3,4), (4,10)}

(10,8)
(10,6)
(4,5)
(2,10)
(1,11)
(1,10)

• • •

Complexity

Elementary operations: Pop and Push

Number of PUSH: $\sum_{v\in V} d(v) = 2m$

Number of POP: $\sum_{v\in V} d(v) = 2m$

O(m)

DFS Algorithm – Recursive version

```
DFS(v)
Mark v visited
 ∀w ∈ Adjacent(v)
         if w not visited
                  visit w
                  DFS(w)
```

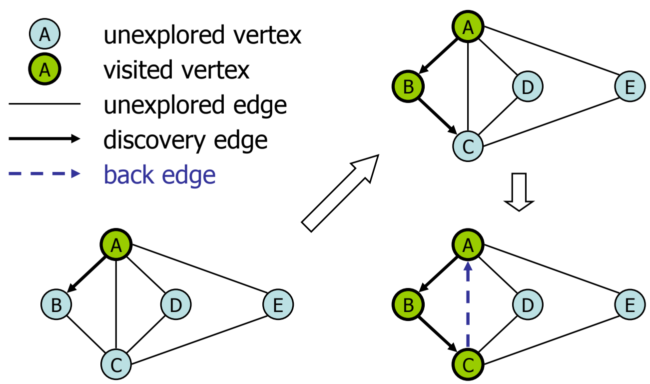

DFS Again - More Detail…

- The algorithm uses a mechanism for setting and getting “labels” of vertices and edges

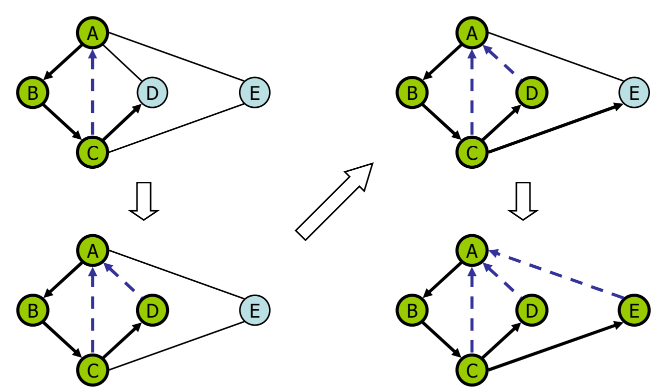
```
Algorithm DFS(G)
    Input graph G
    Output labeling of the edges of G
        as discovery edges and
        back edges
  for all u ∈ G.vertices()
    setLabel(u, UNEXPLORED)
  for all e ∈ G.edges()
    setLabel(e, UNEXPLORED)
  for all v ∈ G.vertices()
    if getLabel(v) = UNEXPLORED
        DFS(G, v)
```

```
Algorithm DFS(G, v)
  Input graph G and a start vertex v of G
  Output labeling of the edges of G
      in the connected component of v
      as discovery edges and back edges
  setLabel(v, VISITED)
  for all e ∈ G.incidentEdges(v)
    if getLabel(e) = UNEXPLORED
      w ← opposite(v,e)
      if getLabel(w) = UNEXPLORED
        setLabel(e, DISCOVERY)
        DFS(G, w)
      else
        setLabel(e, BACK)
```


Example

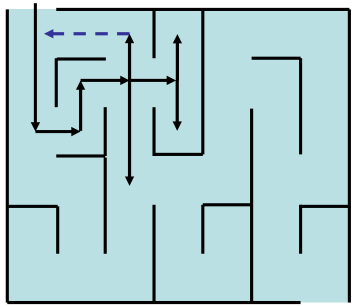


Example (cont.)



DFS and Maze Traversal

- The DFS algorithm is similar to a classic strategy for exploring a maze
  - We mark each intersection, corner and dead end (vertex) visited
  - We mark each corridor (edge) traversed
  - We keep track of the path back to the entrance (start vertex) by means of a rope (recursion stack)



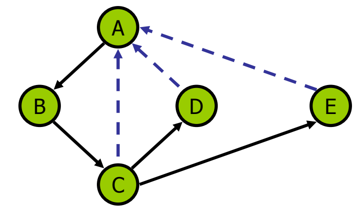
Properties of DFS

Property 1

DFS(*G*, *v*) visits all the vertices and edges in the connected component of *v*

Property 2

The discovery edges labeled by *DFS*(*G*, *v*) form a spanning tree of the connected component of *v*



Analysis of DFS + labeling

- Setting/getting a vertex/edge label takes $O(1)$ time
- Each vertex is labeled twice
  - once as UNEXPLORED
  - once as VISITED

2n

- Each edge is labeled twice
  - once as UNEXPLORED
  - once as DISCOVERY or BACK

2m

- Method incidentEdges is called once for each vertex

$$\sum_{v \in V} d(v) = 2m$$

- DFS runs in $O(n + m)$ time provided the graph is represented by the adjacency list structure

$$O(n + m) = O(m)$$

Conclusion

If we represent the graph with an adjacency list

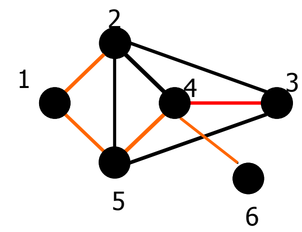
Complexity of DFS is O(m)

WORST CASE: m = O(n^2), when …

Question: Could we do it in less than O(m) ?

Path Finding

- We can specialize the DFS algorithm to find a path between two given vertices u and z using the template method pattern
- We call $DFS(G, u)$ with u as the start vertex
- We use a stack S to keep track of the path between the start vertex and the current vertex
- As soon as destination vertex z is encountered, we return the path as the contents of the stack



2 -- 6

2 1 5 4 ~~3~~ 6

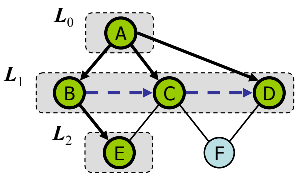

```
Algorithm pathDFS(G, v, z)
  setLabel(v, VISITED)
  S.push(v)
  if v = z
    return S.elements()
  for all e ∈ G.incidentEdges(v)
    if getLabel(e) = UNEXPLORED
      w ← opposite(v,e)
      if getLabel(w) = UNEXPLORED
        setLabel(e, DISCOVERY)
        S.push(e)
        pathDFS(G, w, z)
        S.pop(e)
      else
        setLabel(e, BACK)
  S.pop(v)
```


Cycle Finding

- We can specialize the DFS algorithm to find a simple cycle using the template method pattern
- We use a stack S to keep track of the path between the start vertex and the current vertex
- As soon as a back edge (v, w) is encountered, we return the cycle as the portion of the stack from the top to vertex w

```
Algorithm cycleDFS(G, v, z)
   setLabel(v, VISITED)
   S.push(v)
   for all  e ∈ G.incidentEdges(v)
      if  getLabel(e) = UNEXPLORED
         w ← opposite(v,e)
         S.push(e)
         if getLabel(w) = UNEXPLORED
            setLabel(e, DISCOVERY)
            cycleDFS(G, w, z)
            S.pop(e)
         else
            T ← new empty stack
            repeat
               o ← S.pop()
               T.push(o)
            until o = w
            return T.elements()
   S.pop(v)
```


Breadth-First Search



Outline and Reading

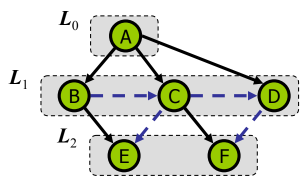
- Breadth-first search (§12.3.3)
  - Algorithm
  - Example
  - Properties
  - Analysis
  - Applications
- DFS vs. BFS
  - Comparison of applications
  - Comparison of edge labels

Breadth-First Search

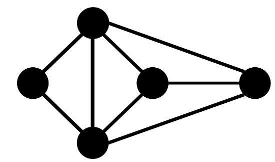
Breadth-First Search is a graph traversal technique that:

- on a graph with n vertices and m edges, takes $O(n + m)$ time
- can be further extended to solve other graph problems
  - Find and report a path with the minimum number of edges between two given vertices
  - Find a simple cycle, if there is one

The idea:
Visit a vertex and then visit all unvisited vertices that are adjacent to it before visiting a vertex which is 2 away from it.

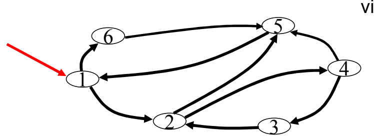


level by level



Simple Breadth-First Search with a Queue

T is the set of visited edges



Visited: {1}
T = ϕ

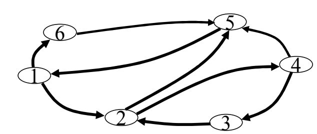
to visit: {(1,2), (1,6)}

(1,2) - Is 2 visited?
Visited: {1,2}
T = {(1,2)}

to visit: {(1,6), (2,4), (2,5)}

(1,6) - Is 6 visited?
Visited: {1,2,6}
T = {(1,2), (1,6)}

to visit: {(2,4), (2,5), (6,5)}



(2,4) - Is 4 visited?
Visited: {1,2,6,4}
T = {(1,2), (1,6), (2,4)}

to visit: {(2,5), (6,5), (4,5), (4,3)}

(2,5) - Is 5 visited ?
Visited: {1,2,6,4,5}
T = {(1,2), (1,6), (2,4), (2,5)}

to visit: {(6,5), (4,5), (4,3)}

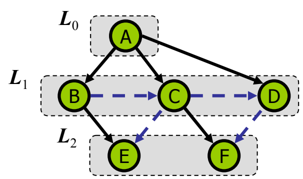
(6,5) - 5? already visited!

(4,5) - 5? already visited!

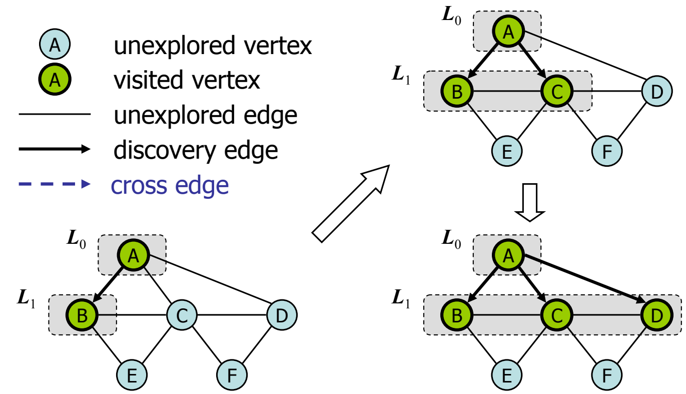
(4,3) - 3?

BFS with labeling

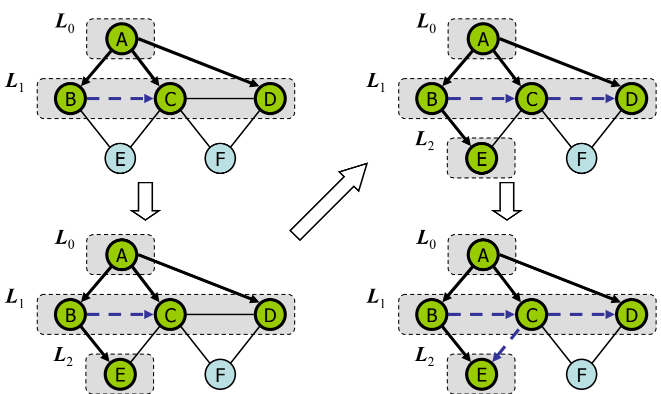
Using a sequence for each level



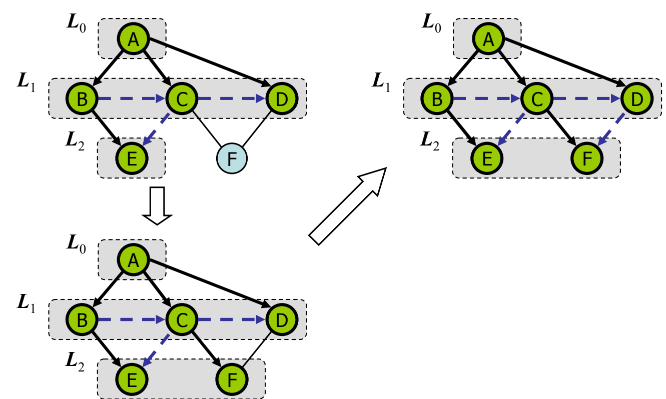
Example



Example (cont.)



Example (cont.)



BFS Again - more details

- The algorithm uses a mechanism for setting and getting "labels" of vertices and edges

```
Algorithm BFS(G)
   Input graph G
   Output labeling of the edges
      and partition of the
      vertices of G
 for all u ∈ G.vertices()
  setLabel(u, UNEXPLORED)
 for all e ∈ G.edges()
  setLabel(e, UNEXPLORED)
 for all v ∈ G.vertices()
  if getLabel(v) = UNEXPLORED
     BFS(G, v)
```

```
Algorithm BFS(G, s)
  L0 ← new empty sequence
  L0.insertLast(s)
  setLabel(s, VISITED)
  i ← 0
  while ! Li.isEmpty()
    Li+1 ← new empty sequence
    for all v ∈ Li.elements()
      for all e ∈ G.incidentEdges(v)
        if getLabel(e) = UNEXPLORED
          w ← opposite(v,e)
          if getLabel(w) = UNEXPLORED
            setLabel(e, DISCOVERY)
            setLabel(w, VISITED)
            Li+1.insertLast(w)
          else
            setLabel(e, CROSS)
    i ← i +1
```


Properties

Notation

G_s: connected component of s

Property 1

$BFS(G, s)$ visits all the vertices and edges of G_s

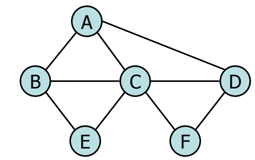
Property 2

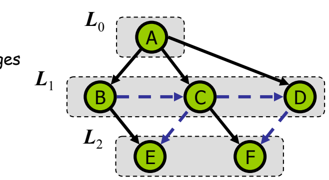
The discovery edges labeled by $BFS(G, s)$ form a spanning tree T_s of G_s

Property 3

For each vertex v in L_i

- The path of T_s from s to v has i edges
- Every path from s to v in G_s has at least i edges





Analysis

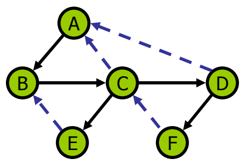
- Setting/getting a vertex/edge label takes $O(1)$ time
- Each vertex is labeled twice
  - once as UNEXPLORED
  - once as VISITED
- Each edge is labeled twice
  - once as UNEXPLORED
  - once as DISCOVERY or CROSS
- Each vertex is inserted once into a sequence L_i
- Method incidentEdges is called once for each vertex
- BFS runs in $O(n + m)$ time provided the graph is represented by the adjacency list structure
  - Recall that $\Sigma_v \deg(v) = 2m$

Applications

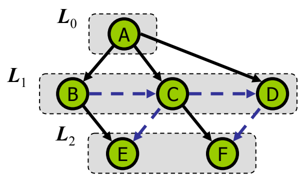
- Using the template method pattern, we can specialize the BFS traversal of a graph G to solve the following problems in $O(n + m)$ time
  - Compute the connected components of G
  - Compute a spanning forest of G
  - Find a simple cycle in G, or report that G is a forest
  - Given two vertices of G, find a path in G between them with the minimum number of edges, or report that no such path exists

DFS vs. BFS

Applications	DFS	BFS
Spanning forest, connected components, paths, cycles	√	√
Shortest paths		√
Biconnected components	√	



DFS

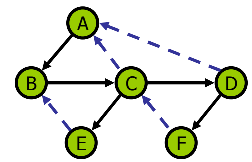


BFS

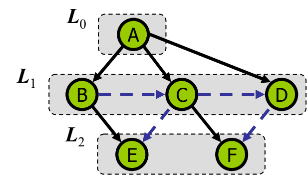
DFS vs. BFS (cont.)

Back edge (v,w)
- w is an ancestor of v in the tree of discovery edges

Cross edge (v,w)
- w is in the same level as v or in the next level in the tree of discovery edges



DFS



BFS