

Introduction Lexicographic Order Maximal/minimal, greatest/least elements Lattices Topological Sorting





There is a total ordering for numbers: we can compare any two numbers and decide which is "larger or equal than".

What about comparing strings? Or sets? Or formulas?

We can come up with a scheme to compare all strings

• lexicographic ordering, will talk about it shortly

However, quite often what we have is that some objects are genuinely incomparable, while for others we can say that **a** is less or equal then **b**.

•In an organization: A being the boss of B

•For sets: A being the subset of B

Captured in the notion of **partial ordering** 



**Definition:** Let **S** be a set and **R** be a binary relation on **S**. If **R** is *reflexive*, *antisymmetric* and *transitive*, then we say that **R** is a partial ordering of **S**. The set **S** together with the partial ordering **R** is called a poset (partially ordered set) and denoted by (**S**, **R**). Members of **S** are called elements of the poset.

We will often use infix notation a R b to denote R(a,b) = T

**Example 1:** S is a set of sets, and the partial order is a set inclusion  $\subseteq$  **Example 2:** S is a set of integers, and R is 'divides':  $aRb \equiv a|b$  **Example 3:** S is a set of formulae, aRb iff a is a sub-formula of b **Example 4:** S is the set of tasks to be scheduled, aRb means "a must finish before b can run"





## **Partial Ordering - Comparability**

We can have the same set S with different relation R:

**Example 5: S** is a set of strings, aRb iff a is a prefix of b

Example 6: S is a set of strings, aRb iff a is a substring of b

**Example 7:** S is a set of formulae, aRb iff  $(b \rightarrow a)$ 

**Example 8:** S is a set of integers, R is the normal  $\leq$ 

The last one is different then previous ones. How?

• every two integers are comparable

**Definition:** The elements **a** and **b** of the poset (S,R) are comparable iff **aRb** or **bRa**. Otherwise, they are called **incomparable**.

**Definition:** If (S,R) is a poset and every two elements of S are comparable, we say that S is totally ordered (or linearly ordered) set and R is called a total (or linear) order.

# Partial Ordering – Lexicographic Order



Lekicographic order - a commonly used technique to construct total order on the Cartesian product of totally ordered sets

- let  $(A_1, R_1)$  and  $(A_2, R_2)$  be totally ordered sets
- $(A_1 \times A_2)$  is a Cartesian product of  $A_1$  and  $A_2$ , essentially a set of the ordered pairs (a,b) such that  $a \in A_1$  and  $b \in A_2$

How can we define a total order R on  $(A_1 \times A_2)$ ?

•  $(a_1, a_2) R (b_1, b_2) iff (a_1 R_1 b_1) or (a_1 = b_1) and (a_2 R_2 b_2)$ 

Straightforwardly extends from 2-tuples to n-tuples

OK, it works for tuples, can we use it for strings?

- works nicely for strings of equal lengths
- what to do if two strings a and b are of different lengths n<m?</p>
  - compare a with the prefix of b of length n
  - if they are equal, then aRb





- How to visualize posets?
- represent elements by nodes
- draw a directed edge from a to b iff aRb

Natural, but messy, with lots of redundant information

#### Hasse Diagrams:

- use the above approach, but delete the edge that must obviously be there by the definition (i.e. the redundant information because of the transitivity and reflexivity of the relation)
- drop the self-cycles
- drop the bridging edges that follow from transitivity
- draw the 'larger' elements higher and drop the arrows
- what is left is the Hasse Diagram of the poset



the poset {1, 2, 3, 4, 6, 8, 12} with the divisibility relation







Not all elements in the poset are equal

- $a \in S$  is maximal in (S,R), iff  $\neg(\exists b: (aRb \land a \neq b))$
- $a \in S$  is minimal in (S,R), iff  $\neg(\exists b: (bRa \land a \neq b))$
- $a \in S$  is a greatest element in (S,R), iff  $\forall b: bRa$
- $a \in S$  is a least element in (S,R), iff  $\forall b: aRb$

Is it true that every maximal element is also a greatest element? Is it true that every greatest element is also a maximal element? Is it true that a poset might have several maximal elements? Is it true that a poset might have several least elements?





### **Partial Ordering – Special Elements**



Maximal elements: 8 and 12 Minimal elements: 1 Greatest element: NONE Least element: 1





Let A be a subset of S, where (S,R) is a poset

We say that **b** is an **upper bound** of the set **A** iff  $\forall a \in A$ , **aRb** 

We say that **b** is a **lower bound** of the set **A** iff  $\forall a \in A$ , **bRa** 

We say that **c** is the least upper bound (denoted by **lub(A)**) bound of the set **A** iff for every upper bound **b** of **A** holds **cRb** 

We say that **c** is the greatest lower bound (denoted by **glb(A)** of the set **A** iff for every lower bound **b** of **A** holds **bRc** 









8 is an upper bound for the set {4, 6} 12 is an upper bound for the set {4, 6} But there are no least upper bound for the set {4, 6}

1 and 2 are both lower bounds for the set {4, 6}2 is the greatest lower bounds for the set {4, 6}

Note that for the poset (Z+, |), glb(A) coincides with GCD(A) and lub(A) coincides with LCM(A).



# **Partial Ordering – Lattices**



The posets for which every pair of elements have a **glb/lub** have many special properties and applications and are of special interest

- by induction, we can prove that every finite non-empty subset
   A of S has glb/lub, i.e. considering only pairs is not limiting
- Such posets are called **lattices**
- Is (Z<sup>+</sup>, ) a lattice?
- Is ({1,2,4,8,12,24}, |) a lattice?
- Is  $(Z^+, \leq)$  a lattice?

Let S be a set with n elements. Is the sets of all subsets of S ordered by inclusion "power set of S" a lattice?





Given a set of tasks and dependences between them, find a schedule (for one computer) to execute them all.

What we want is a **compatible** total order **R'**:

∀a,b∈S: if aRb then also aR'b

also called topological sorting of (S,R)

Actually, R' does not need to be total order to say that it is compatible with R

 but to say that R' is topological sorting, we require that R' is total



Constructing a Topological Sorting:

Input: A finite poset (S,R)

```
k = 1;
while (S ≠Ø) {
    a<sub>k</sub> = a minimal element of S; // such element must exist
    S = S - {ak};
    k++;
}
```

**Output:** The sequence  $\{a_1, a_2, \dots, a_n\}$  defining the total order R' of S.





How do we prove a statement P(x) for every element x of the poset? Quite often, the following approach is very natural:

- Basis step: Directly prove P(x) for every minimal element x
- Induction step: Prove that if P(y) for all y such that yRx then also P(x)

Does this work for all posets?

- no, as there are posets that have infinite chains of smaller and smaller elements (e.g.  $(Z, \leq)$ )
- it works for well-founded posets
  - every non-empty set has a minimal element

• differs from the **well-ordered sets** by not requiring that the order relation be total