

An aerial photograph of a city, likely Ottawa, showing a river flowing through a green park area with several bridges. In the background, there are modern high-rise buildings under a blue sky with scattered clouds. A dark semi-transparent banner is overlaid on the middle of the image, containing the title and author information.

# ELG4177 - DIGITAL SIGNAL PROCESSING

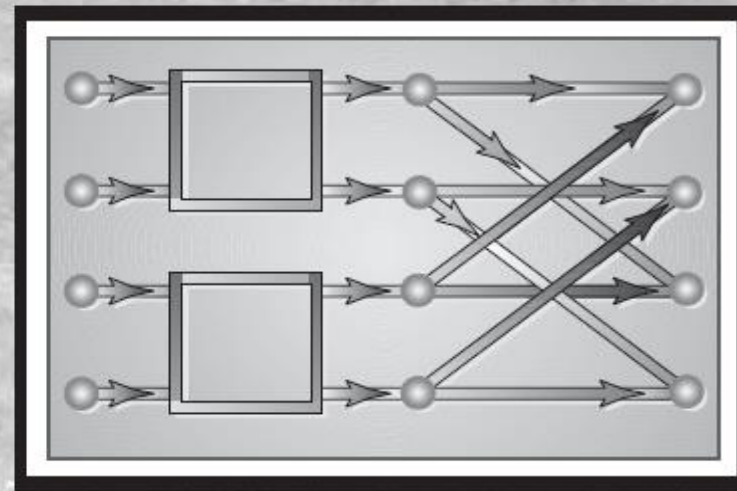
## Tutorial 10

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## 9

Computation  
of the Discrete  
Fourier Transform

Problems: 9.3, 9.5, 9.6, 9.7, 9.13, 9.15, 9.16, 9.18, 9.19, 9.23, 9.26, 9.28

- 9.3. Suppose that you time-reverse and delay a real-valued 32-point sequence  $x[n]$  to obtain  $x_1[n] = x[32-n]$ . If  $x_1[n]$  is used as the input for the system in Figure P9.4, find an expression for  $y[32]$  in terms of  $X(e^{j\omega})$ , the DTFT of the original sequence  $x[n]$ .

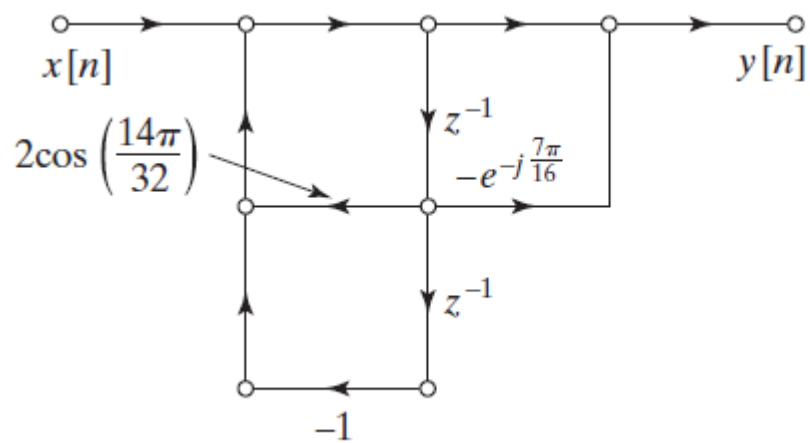


Figure P9.4

First, we derive a relationship between the  $X_1(e^{j\omega})$  and  $X(e^{j\omega})$  using the shift and time reversal properties of the DTFT.

$$\begin{aligned}x_1[n] &= x[32 - n] \\X_1(e^{j\omega}) &= X(e^{-j\omega})e^{-j32\omega}\end{aligned}$$

Looking at the figure we see that calculating  $y[32]$  is just an application of the Goertzel algorithm with  $k = 7$  and  $N = 32$ . Therefore,

$$\begin{aligned}y[32] &= X_1[7] \\&= X_1(e^{j\omega})\big|_{\omega=\frac{2\pi 7}{32}} \\&= X(e^{-j\omega})e^{-j\omega 32}\big|_{\omega=\frac{7\pi}{16}} \\&= X(e^{-j\frac{7\pi}{16}})e^{-j(\frac{7\pi}{16})32} \\&= X(e^{-j\frac{7\pi}{16}})\end{aligned}$$

9.5. Consider the signal flow graph in Figure P9.5. Suppose that the input to the system  $x[n]$  is an 8-point sequence. Choose the values of  $a$  and  $b$  such that  $y[8] = X(e^{j6\pi/8})$ .

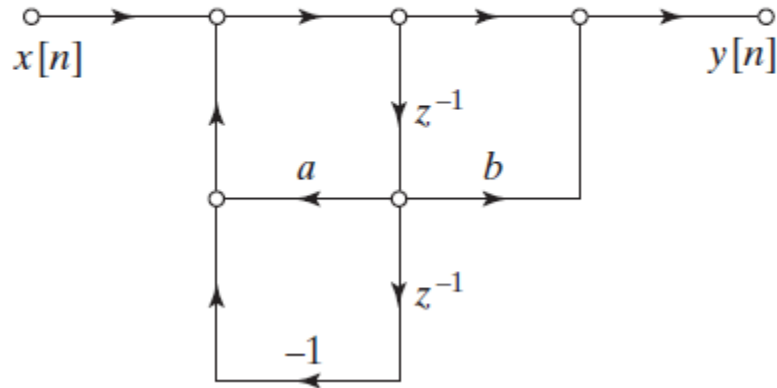


Figure P9.5

.  $X(e^{j6\pi/8})$  corresponds to the  $k = 3$  index of a length  $N = 8$  DFT. Using the flow graph of the second-order recursive system for Goertzel's algorithm,

$$\begin{aligned}
 a &= 2 \cos\left(\frac{2\pi k}{N}\right) \\
 &= 2 \cos\left(\frac{2\pi(3)}{8}\right) \\
 &= -\sqrt{2} \\
 b &= -W_N^k \\
 &= -e^{-j6\pi/8} \\
 &= \frac{1+j}{\sqrt{2}}
 \end{aligned}$$

9.6. Figure P9.6 shows the graph representation of a decimation-in-time FFT algorithm for  $N = 8$ . The heavy line shows a path from sample  $x[7]$  to DFT sample  $X[2]$ .

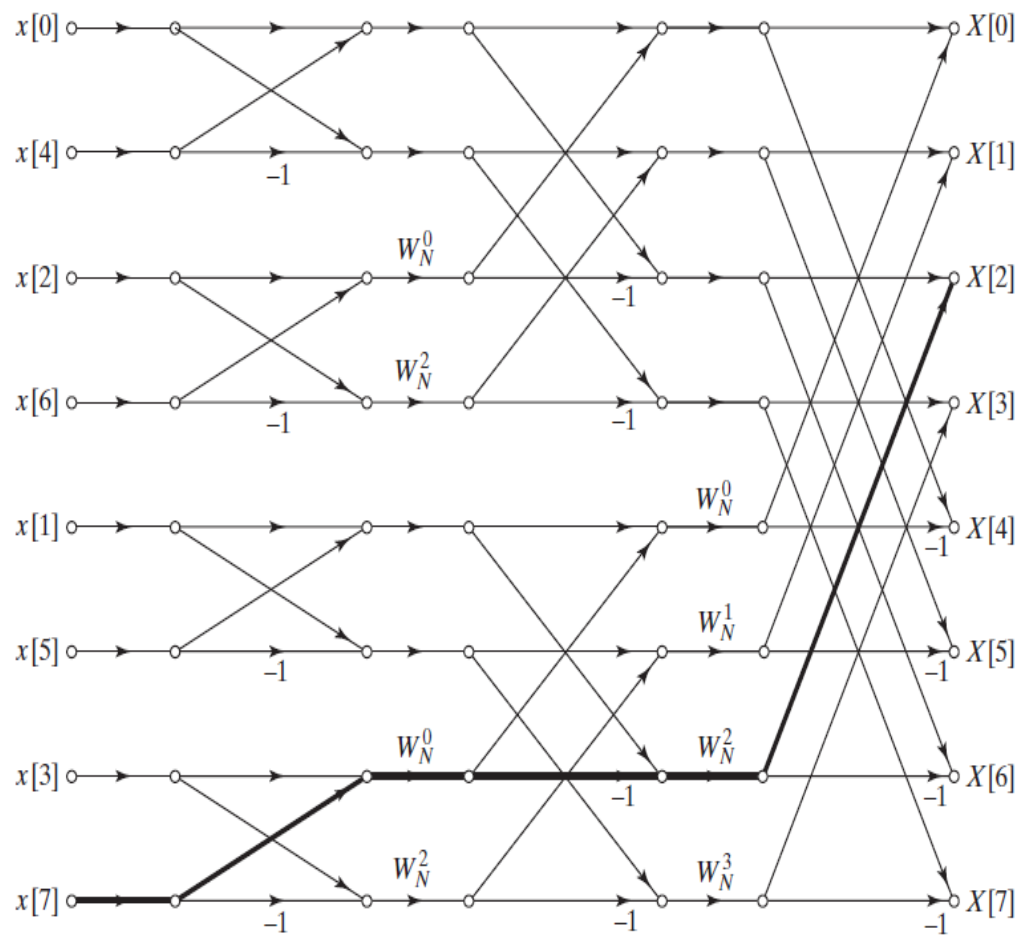


Figure P9.6

- (a) What is the “gain” along the path that is emphasized in Figure P9.6?
- (b) How many other paths in the flow graph begin at  $x[7]$  and end at  $X[2]$ ? Is this true in general? That is, how many paths are there between each input sample and each output sample?
- (c) Now consider the DFT sample  $X[2]$ . By tracing paths in the flow graph of Figure P9.6, show that each input sample contributes the proper amount to the output DFT sample; i.e., verify that

$$X[2] = \sum_{n=0}^{N-1} x[n]e^{-j(2\pi/N)2n}.$$

- (a) The "gain" along the emphasized path is  $-W_N^2$ .
- (b) In general, there is only one path between each input sample and each output sample.
- (c)  $x[0]$  to  $X[2]$ : The gain is 1.  
 $x[1]$  to  $X[2]$ : The gain is  $W_N^2$ .  
 $x[2]$  to  $X[2]$ : The gain is  $-W_N^0 = -1$ .  
 $x[3]$  to  $X[2]$ : The gain is  $-W_N^0 W_N^2 = -W_N^2$ .  
 $x[4]$  to  $X[2]$ : The gain is  $W_N^0 = 1$ .  
 $x[5]$  to  $X[2]$ : The gain is  $W_N^0 W_N^2 = W_N^2$ .  
 $x[6]$  to  $X[2]$ : The gain is  $-W_N^0 W_N^0 = -1$ .  
 $x[7]$  to  $X[2]$ : The gain is  $-W_N^0 W_N^0 W_N^2 = -W_N^2$ , as in Part (a).

Now

$$\begin{aligned}
 X[2] &= \sum_{n=0}^7 x[n] W_8^{2n} \\
 &= x[0] + x[1] W_8^2 + x[2] W_8^4 + x[3] W_8^6 + x[4] W_8^8 + x[5] W_8^{10} + x[6] W_8^{12} \\
 &\quad + x[7] W_8^{14} \\
 &= x[0] + x[1] W_8^2 + x[2](-1) + x[3](-W_8^2) + x[4](1) + x[5] W_8^2 \\
 &\quad + x[6](-1) + x[7](-W_8^2)
 \end{aligned}$$

Each input sample contributes the proper amount to the output DFT sample.



9.7. Figure P9.7 shows the flow graph for an 8-point decimation-in-time FFT algorithm. Let  $x[n]$  be the sequence whose DFT is  $X[k]$ . In the flow graph,  $A[\cdot]$ ,  $B[\cdot]$ ,  $C[\cdot]$ , and  $D[\cdot]$  represent separate arrays that are indexed consecutively in the same order as the indicated nodes.

- (a) Specify how the elements of the sequence  $x[n]$  should be placed in the array  $A[r]$ ,  $r = 0, 1, \dots, 7$ . Also, specify how the elements of the DFT sequence should be extracted from the array  $D[r]$ ,  $r = 0, 1, \dots, 7$ .
- (b) Without determining the values in the intermediate arrays,  $B[\cdot]$  and  $C[\cdot]$ , determine and sketch the array sequence  $D[r]$ ,  $r = 0, 1, \dots, 7$ , if the input sequence is  $x[n] = (-W_N)^n$ ,  $n = 0, 1, \dots, 7$ .
- (c) Determine and sketch the sequence  $C[r]$ ,  $r = 0, 1, \dots, 7$ , if the output Fourier transform is  $X[k] = 1$ ,  $k = 0, 1, \dots, 7$ .

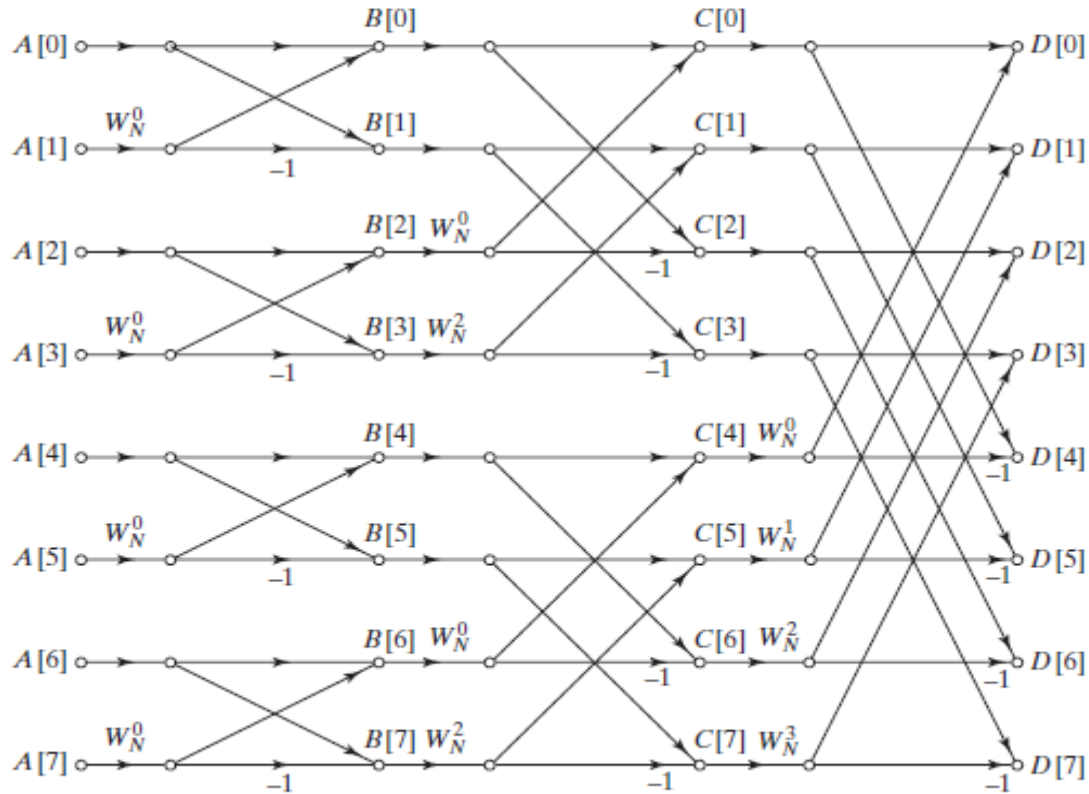


Figure P9.7



(a) The input should be placed into  $A[r]$  in bit-reversed order.

$$\begin{aligned}
 A[0] &= x[0] \\
 A[1] &= x[4] \\
 A[2] &= x[2] \\
 A[3] &= x[6] \\
 A[4] &= x[1] \\
 A[5] &= x[5] \\
 A[6] &= x[3] \\
 A[7] &= x[7]
 \end{aligned}$$

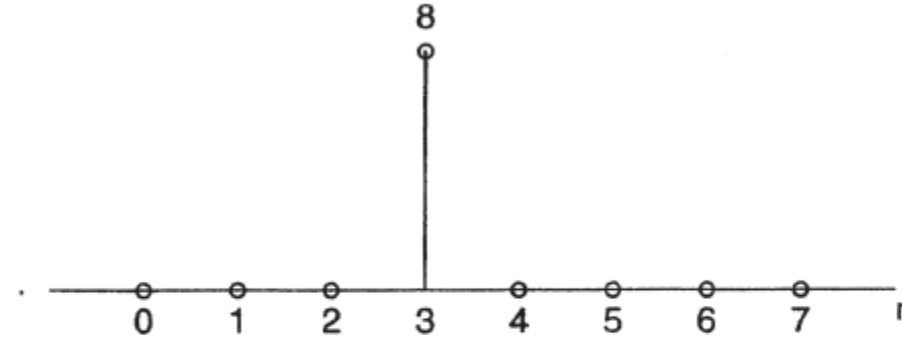
The output should then be extracted from  $D[r]$  in sequential order.

$$X[k] = D[k], \quad k = 0, \dots, 7$$

(b) First, we find the DFT of  $(-W_N)^n$  for  $N = 8$ .

$$\begin{aligned}
 X[k] &= \sum_{n=0}^7 (-W_8)^n W_8^{nk} \\
 &= \sum_{n=0}^7 (-1)^n W_8^n W_8^{nk} \\
 &= \sum_{n=0}^7 (W_8^{-4})^n W_8^n W_8^{nk} \\
 &= \sum_{n=0}^7 W_8^{n(k-3)} \\
 &= \frac{1 - W_8^{k-3}}{1 - W_8^{k-3}} \\
 &= 8\delta[k-3]
 \end{aligned}$$

A sketch of  $D[r]$  was provided below.



(c) First, the array  $D[r]$  is expressed in terms of  $C[r]$ .

$$\begin{aligned}
 D[0] &= C[0] + C[4] \\
 D[1] &= C[1] + C[5]W_8^1 \\
 D[2] &= C[2] + C[6]W_8^2 \\
 D[3] &= C[3] + C[7]W_8^3 \\
 D[4] &= C[0] - C[4] \\
 D[5] &= C[1] - C[5]W_8^1 \\
 D[6] &= C[2] - C[6]W_8^2 \\
 D[7] &= C[3] - C[7]W_8^3
 \end{aligned}$$

Solving this system of equations for  $C[r]$  gives

$$C[0] = (D[0] + D[4])/2$$

$$C[1] = (D[1] + D[5])/2$$

$$C[2] = (D[2] + D[6])/2$$

$$C[3] = (D[3] + D[7])/2$$

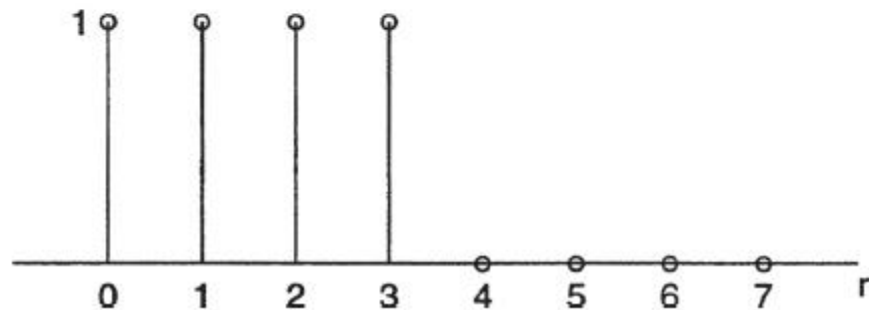
$$C[4] = (D[0] - D[4])/2$$

$$C[5] = (D[1] - D[5])W_8^{-1}/2$$

$$C[6] = (D[2] - D[6])W_8^{-2}/2$$

$$C[7] = (D[3] - D[7])W_8^{-3}/2$$

for  $r = 0, 1, \dots, 7$ . A sketch of  $C[r]$  is provided below.



**9.13.** Assume that you wish to sort a sequence  $x[n]$  of length  $N = 16$  into bit-reversed order for input to an FFT algorithm. Give the new sample order for the bit-reversed sequence.

Reversing the bits (denoted by  $\rightarrow$ ) gives

0	=	0000	$\rightarrow$	0000	=	0
1	=	0001	$\rightarrow$	1000	=	8
2	=	0010	$\rightarrow$	0100	=	4
3	=	0011	$\rightarrow$	1100	=	12
4	=	0100	$\rightarrow$	0010	=	2
5	=	0101	$\rightarrow$	1010	=	10
6	=	0110	$\rightarrow$	0110	=	6
7	=	0111	$\rightarrow$	1110	=	14
8	=	1000	$\rightarrow$	0001	=	1
9	=	1001	$\rightarrow$	1001	=	9
10	=	1010	$\rightarrow$	0101	=	5
11	=	1011	$\rightarrow$	1101	=	13
12	=	1100	$\rightarrow$	0011	=	3
13	=	1101	$\rightarrow$	1011	=	11
14	=	1110	$\rightarrow$	0111	=	7
15	=	1111	$\rightarrow$	1111	=	15

The new sample order is 0, 8, 4, 12, 2, 10, 6, 14, 1, 9, 5, 13, 3, 11, 7, 15.

- 9.15. The butterfly in Figure P9.15 was taken from a decimation-in-frequency FFT with  $N = 16$ , where the input sequence was arranged in normal order. Note that a 16-point FFT will have four stages, indexed  $m = 1, \dots, 4$ . Which of the four stages have butterflies of this form? Justify your answer.

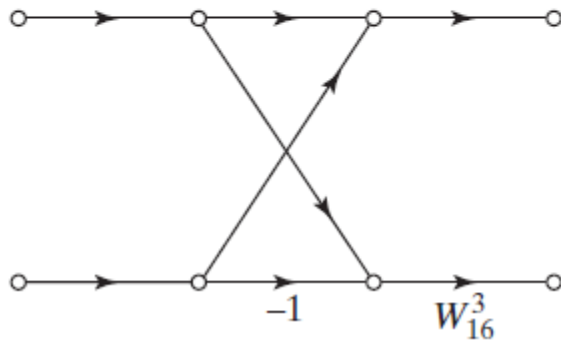


Figure P9.15

### Solution

Only the  $m = 1$  stage will have this form. No other stage of a  $N = 16$  radix-2 decimation-in-frequency FFT will have a  $W_{16}$  term raised to an odd power.



9.16. The butterfly in Figure P9.16 was taken from a decimation-in-time FFT with  $N = 16$ . Assume that the four stages of the signal flow graph are indexed by  $m = 1, \dots, 4$ . What are the possible values of  $r$  for each of the four stages?

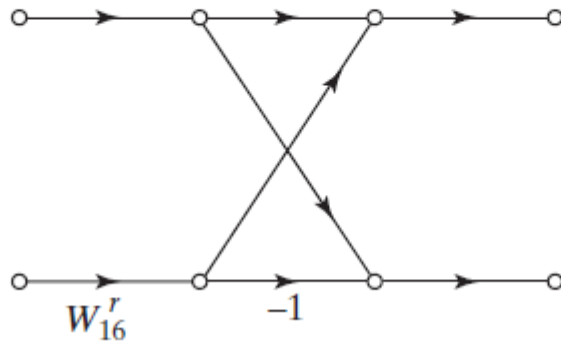


Figure P9.16

Solution

The possible values of  $r$  for each of the four stages are

$$\begin{aligned}
 m = 1, & \quad r = 0 \\
 m = 2, & \quad r = 0, 4 \\
 m = 3, & \quad r = 0, 2, 4, 6 \\
 m = 4, & \quad r = 0, 1, 2, 3, 4, 5, 6, 7
 \end{aligned}$$

9.18. The butterfly in Figure P9.18 was taken from a decimation-in-time FFT with  $N = 16$ . Assume that the four stages of the signal flow graph are indexed by  $m = 1, \dots, 4$ . Which of the four stages have butterflies of this form?

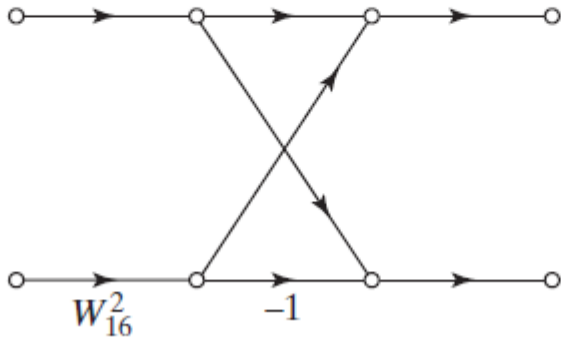


Figure P9.18

Solution

The possible values for  $r$  for each of the four stages are

- $m = 1, \quad r = 0$
- $m = 2, \quad r = 0, 4$
- $m = 3, \quad r = 0, 2, 4, 6$
- $m = 4, \quad r = 0, 1, 2, 3, 4, 5, 6, 7$

where  $W_N^r$  is the twiddle factor for each stage. Since the particular butterfly shown has  $r = 2$ , the stages which have this butterfly are

$$m = 3, 4$$

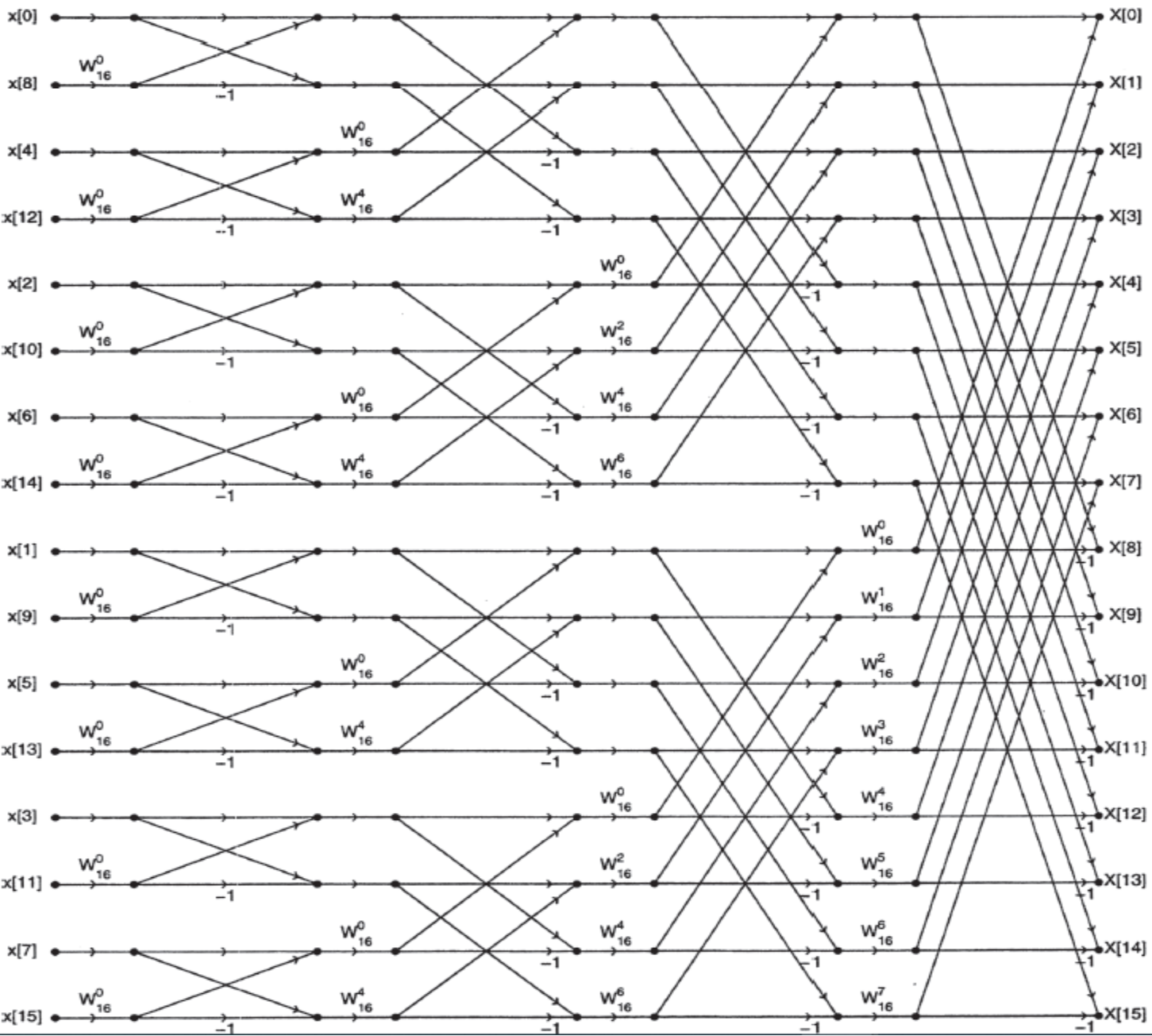
**9.19.** Suppose you are told that an  $N = 32$  FFT algorithm has a “twiddle” factor of  $W_{32}^2$  for one of the butterflies in its fifth (last) stage. Is the FFT a decimation-in-time or decimation-in-frequency algorithm?

### Solution

The FFT is a decimation-in-time algorithm, since the decimation-in-frequency algorithm has only  $W_{32}^0$  terms in the last stage.

- 9.23.** Construct a flow graph for a 16-point radix-2 decimation-in-time FFT algorithm. Label all multipliers in terms of powers of  $W_{16}$ , and also label any branch transmittances that are equal to  $-1$ . Label the input and output nodes with the appropriate values of the input and DFT sequences, respectively.





**9.26.** We are given a finite-length sequence  $x[n]$  of length 627 (i.e.,  $x[n] = 0$  for  $n < 0$  and  $n > 626$ ), and we have available a program that will compute the DFT of a sequence of any length  $N = 2^v$ .

For the given sequence, we want to compute samples of the DTFT at frequencies

$$\omega_k = \frac{2\pi}{627} + \frac{2\pi k}{256}, \quad k = 0, 1, \dots, 255.$$

Specify how to obtain a new sequence  $y[n]$  from  $x[n]$  such that the desired frequency samples can be obtained by applying the available FFT program to  $y[n]$  with  $v$  as small as possible.

## Solution

Let

$$y[n] = e^{-j2\pi n/627} x[n]$$

Then

$$Y(e^{j\omega}) = X(e^{j(\omega + \frac{2\pi}{627})})$$

Let  $y'[n] = \sum_{m=-\infty}^{\infty} y[n + 256m]$ ,  $0 \leq n \leq 255$ , and let  $Y'[k]$  be the 256 point DFT of  $y'[n]$ . Then

$$Y'[k] = X\left(e^{j\left(\frac{2\pi k}{256} + \frac{2\pi}{627}\right)}\right)$$

9.28. You are asked to build a system that computes the DFT of a 4-point sequence

$$x[0], x[1], x[2], x[3].$$

You can purchase any number of computational units at the per-unit cost shown in Table 9.1.

**TABLE 9.1**

Module	Per-Unit Cost
8-point DFT	\$1
8-point IDFT	\$1
adder	\$10
multiplier	\$100

Design a system of the lowest possible cost. Draw the associated block diagram and indicate the system cost.

A. The DFT of the 4-point sequence  $x[0], x[1], x[2], x[3]$  is given by

$$X[k] = X(e^{j\omega}) \Big|_{\omega=\frac{2\pi k}{4}}, \quad k = 0, 1, 2, 3,$$

where

$$X(e^{j\omega}) = \sum_{n=0}^3 x[n] e^{-j\omega n}.$$

The 8-point DFT of the sequence  $x[0], x[1], x[2], x[3], 0, 0, 0, 0$  is given by

$$\hat{X}[k] = X(e^{j\omega}) \Big|_{\omega=\frac{2\pi k}{8}}, \quad k = 0, \dots, 7,$$

where  $X(e^{j\omega})$  is as above. We see that

$$X[k] = \hat{X}[2k], \quad k = 0, 1, 2, 3.$$

The cost of the system is the cost of the 8-point DFT, which is \$1.