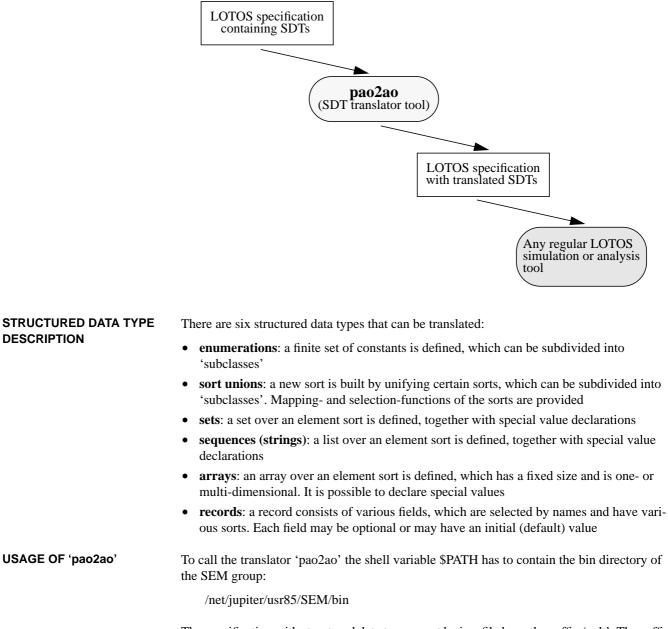
The Structured Data Type Translator 'pao2ao'

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ACT ONE is an abstract data type language that is very flexible, because it allows to define any function indepedently from any application. This advantage is unfortunatly undermined by the fact that abstract data type has to be systematically constructed from scratch, thus requiring sometimes considerable work and thus reducing its usability. Extensions to the language have never taken place due to the objectives of stability defined by the various ISO work group. Experience by various LOTOS community users has shown that some frequently used concepts should be implemented in ACT ONE in a standard way. This has resulted in several proposals [SS91] [UT92]. Both consisted in defining an extension to the language with semantics to translate these extensions into traditional ACT ONE. The following work consisted in writing a translator that would take a LOTOS specification where new Structured Data Types (SDT) are present both as definitions, but also as value expressions in actions, process instances value actualizations and guards.

There were several strategies possible. The decision to build a stand alone translator was based on several factors. Among them is the fact that several LOTOS tools exist, and a translator would have avoided to modify all of the existing tools. Also, the Structured Data type are presently experimental and are not included yet in the standard. The translator is then considered as allowing a user to reason in a higher level of abstraction independent from ACT ONE itself. In this case the resulting ACT ONE translation is considered as strictly an executable version similar to assembly language in traditional computer languages. (fig.1)

FIGURE 1. Tool chain with SDT translation tool 'pao2ao'



The specification with structured data types must be in a file have the suffix '.sdt'. The suffix is different or there is no suffix, then the tool will look for a file 'name.sdt'. The tool generates a file 'name.lot'. **Caution**: the tool does not check whether there exists already a file with that name.

A help message, saying how the tool is called and which options are possible, can be obtained by using the '-h' option:

```
> pao2ao -hef
```

usage:	sdt2ao inputfile[.sdt] [-o outputfile] [-g]				
	where [-g]: generation of visualisation annotations				
	and Demon graphic template program				
or:	sdt2ao -h				
	which prints this message				

A successfull translation produces the following output. (Please notice that the file containing the specification with structured data types needs to have the suffix '.sdt'. The suffix can be omitted):

> pao2ao test

Gesellschaft fuer Mathematik und Datenverarbeitung Forschungszentrum fuer Offene Kommunikationssysteme Structured Data Type to Act One Translator

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Translation of file: test.sdt starts !

Parsing and syntax checking start ! Parsing and syntax ckecking were successful !

Unparsing starts ! Output is written to file: 'test.lot'

Unparsing successful !

Translation of Structured Data Types was completely successful !!!

Interpreting Error Reports of the Static Semantic Checker

	The translator 'pao2ao' only checks the correctness of the syntax of a specification. It does not check the static semantic. This has to be done after the translation by using commonly available LOTOS tools, like for instance TOPO that is distributed together with the <i>lite</i> tool. The interpretation of the errors, reported by these tools, is sometimes a bit tricky because they point to the translated specification.
	In the following we will list some typical error messages and give guideline how to interpret them. Of cause the list can not be complete and there are many, many other cases possible. We have used the TOPO compiler for the static semantic check.
	In most cases, the best strategy is look only at the errors reported at the first 1 or 2 lines and try to correct them. The following are probably inferences of the first errors.
TYPE NOT FOUND	type "Set" not found at library type "NaturalNumber" not found at library type "String" not found at library
	Library must be changed. The structured data types from the library. The library must pro- vide these three types and order relation on the terms. For lite the mod-is library is necessary (see annex for the required signature of the three types from above).
sorts of value expression are different / undefined operation	test.lot:165: lsa: sorts of value expressions are different test.lot:165: lsa: undefined operation: array test.lot:165: lsa: undefined operation: a
	The operation 'a' is not defined.

test.lot:82: lsa: sorts of value expressions are different test.lot:82: lsa: undefined operation: array

The number of arguments of the array is wrong.

ina.lot:1012: lsa: undefined operation: overlay

This message indicates that the record name might be wrong. Each record is translated into a set of record fields which is overlayed over a set of initialization values.

Signatures of Translated Structured Data Types

Structured data types use four library data types from the '**mod-is**' library: String, Set, NaturalNumber, Boolean. These types are imported from the library, but only if a type is used, and no type exists in the specification that has the same name. The translator tool also generates two other types:

1. DecimalNaturalNumber: enrichment of 'NaturalNumber' by functions which are mapped on successors of 0 (used for the array indexes)

	11	· ·	-		·	
	1, 2, 3, 4, 5, 6, 7, 8, 9:		->	Nat		
	_ · _:	Nat, Nat	->	Nat		
•	sequence_SDT: enrich	hment of 'String' l	by th	e functio	ns	
	rpush, lpush:	Element, String	->	String	(* add element to right	*)
					(* left end of sequence	*)
	rpop, lpop:	String	->	String	(* right/left sequence with	*)
					(* one element removed	*)
	right, left:	String	->	Element	(* right/left element	*)
	lt:	String, String	->	Bool	(* compares elementwise	*)
					· •	

formal operation:

2.

lt: Element, Element -> FBool

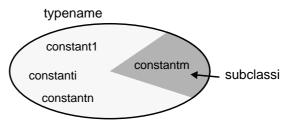
The syntax of the terms of the structured data types that can be used in the specification can be found in the values definitions of the structured data types.

ENUMERATION

The enumeration type is used to introduce a finite set of constants. The constants of an enumeration type can be used directly by their name. The constants are ordered by the sequence in which they are declared. The first one is the least. Constants are used in a specification

- as **basic building elements**, which have no finer structure, like for example in the OSI environment the informations about the kind of connections (e.g. negotiated or not) or the result of service invocations (e.g. ok or failure reason)
- as **abstractions of complex data structures**. In the initial stage of specification evolution, the internal structure of data elements is not important. The constants are later transformed into complex types in implementation directed specifications. Furthermore constants can be used instead of complex data structures to get clearer simulations of process behaviour.

Constants can be grouped into subclasses of a type. For each subclass an additional predicate is specified. These predicates have the form 'is_subclassi(x)', where x is a variable or a term.



Enumeration Scheme	enumtype is	typename			
		{ constant1, constant2,,			
	subclass	subclass1 { constantj,, co		(* OPTIONAL	*)
		subclass2 { constantl,, co	(* OPTIONAL	*)	
		 subclassn { constantm,,	constantm }	(* OPTIONAL	*)
	endtype				
Types	typename				
Library Types	NaturalNumb	per (Import)			
Sorts	typename				
Operations	constant1:		-> typename	e (* constructors	*)
- F	constant1:		-> typename		,
			<i></i>		
	constantn:		-> typename		
	next:	typename	-> typename	e (* successor-operation	*)
	min_typenam	ne:	-> typename	e (* minimum&maximum	1*)
	max_typenan	ne:	-> typename	e (* of enumeration	*)
	is_constant1:	typename	-> Bool	(* predicates	*)
	is_constant2:	typename	-> Bool		
	 is_constantn:	typename	-> Bool		
	eq:	typename, typename	-> Bool	(* equality & order rel's	, *)
	ne:	typename, typename	-> Bool		
	lt:	typename, typename	-> Bool		
	le:	typename, typename	-> Bool		
	gt:	typename, typename	-> Bool		
	ge:	typename, typename	-> Bool		
	is_subclass1:	••	-> Bool	(* subclass predicates	*)
	is_subclass2:	typename	-> Bool	(* constant is in (* subclass	*)
	is_subclassn:	typename	-> Bool		
	Auxiliary Op	erations:			
	h:	typename	-> Nat	(* mapping of constants (* on Nat for order rel's	

UNION

Union types construct a new data type as the union of given data types (called element types). In contrast to the type combinations of standard LOTOS, they generate a new sort that collects the data objects of the unified types. They also provide the necessary operations to manage the union like test the type of an element, select an element and set the value of an element. Usual applications of union types can be found in the field of specifying OSI-Communications, where the protocol services are grouped according to their different functional-ities.

LOTOS does not include hierarchical sorts / classes in its basic mathematical model. An example of such sort hierarchies is the specification of natural numbers as a sub sort of integers. The integer sort is assignment compatible with the natural numbers, which means that a natural value can be assigned to an integer variable. LOTOS can just simulate hierarchical sorts by providing type conversion functions and predicates to test the subsort membership of values. This is done by union types.

Union types could be seen as superclasses of other types. A superclass unifies data objects that share some properties. Within a union of objects, the objects can be grouped into subclasses.

The operations that can be performed on unions are the selection of union elements and the setting of the values of union elements. It is possible to test whether a union component is of a certain type. The conceptual framework provides equality and type testing predicates and order relations on unions.

Union Scheme	is in	constructor2: impor	t_sort1, t_sort2, t_sortn }	(* OPTIONAL	*)
		ubclass2 { constructorl,	, constructork }	(* OPTIONAL	*)
		ubclassn { constructorm, .	, constructorm }	(* OPTIONAL	*)
Types	typename				
Library Types	Boolean (Impo	ort)			
Sorts	typename				
Operations	constructor1: constructor2: constructorn:	sort1 sort2 sortn	-> typename-> typename-> typename	(* constructors: (* map import_sort to (* union_sort	*) *) *)
	is_constructor is_constructor is_constructor	2: typename	-> Bool -> Bool -> Bool	(* predicates	*)
	get_constructo get_constructo get_constructo	r2: typename	-> sort1 -> sort2 -> sortn	(* selectors: (* converts union_sort (* back to import_sort	*) *) *)
	eq: _ne_: _lt_:	typename, typename typename, typename typename, typename	-> Bool -> Bool -> Bool	(* equality (* lexicographical orde	*) er *)
	is_subclass1: is_subclass2:	typename typename	-> Bool -> Bool	(* subclass predicates: (* term has one of the (* sorts described by	*) *) *)
	is_subclassn:	typename	-> Bool	(* the constructor set (* of that subclass) *) *)

SEQUENCE

Sequences are, like arrays, ordered sequences of elements where all elements have the same sort. All components of a sequence can be read but only the outermost component can be modified. The modification of inner components is only possible by first removing one by one all components before the component can be accessed.

Sequences are used in LOTOS specifications, in case the number of components is not fixed from the beginning. An example for sequences in the OSI range are the sequences of acquaintances.

The operations on sequences are append elements to the sequence, select its head or its tail, select the *n*th element of a sequence concatenate two sequences and get the length of a sequence. The conceptual framework provides equality and type testing predicates and order relations on sequences.

Sequence Scheme	seqtype is elements values endtype	typename import_type sort constant1 = (expressioni, constant2 = (expressionk, . constantn = (expressionm,	, expressionl) ;	(* OPTIONAL (* OPTIONAL (* OPTIONAL	*) *) *)
Types	typename				
	••	typename00			
Library Types	sequence_S	DT (Actualized and Renamed)		
Sorts	typename				
Operations	<>: _+_: typename: rpush: lpush: rpop: lpop: right: left: _++:	typename, sort sort, typename sort sort, typename sort, typename typename typename typename typename typename	 -> typename -> sort -> sort -> sort -> typename 	(* constructors (* seq. of only one elem (* add element to right/ (* left end of sequence (* right/left sequence (* with one element (* removed (* right/left element (* seq. concatenation	
	Length: Reverse:	typename	-> Nat -> typename		,
	eq: _ne_: _lt_: constant1: constant2: constantn:	typename, typename typename, typename typename, typename	 -> Bool -> Bool -> Bool -> typename -> typename -> typename 	(* equality (* lexicographical order (* constants for seq (* values	*) r *) *) *)

Sets are collections of elements where all elements have the same sort. Unlike sequences, sets are not ordered.

The usual operations on sets the test of membership for values, the union and intersection of sets and the test if one set is a subset of another one.

Set Scheme		typename				
		import_type				
		sort		• • • • •		*
		$constant1 = \{ expressioni, .$			(* OPTIONAL	*) *)
		$constant2 = \{ expressionk, \}$, expr	ession1 } ;	(* OPTIONAL	*)
		 constantn = { expressionm,	exp	ressionn 11.	(* OPTIONAL	*)
	endtype		, exp	ressionin j ₁ ,	(OF HORAL)
Types	typename					
	Auxiliary Typ	6 5.				
	typename0, ty					
	of permittee, ty	r				
Library Types	Set (Actualize	ed and Renamed)				
Sorts	typename					
Operations	{}:		->	typename	(* constructors	*)
- por anono	Insert:	sort, typename		typename	(* element in set, but) *)
	moort.	sort, typenume	-	., pename	(* only if it is new	*)
					× •	,
	Remove:	sort, typename	->	typename	(* remove element	*)
	IsIn:	sort, typename		Bool	(* tests whether an	*)
	NotIn:	sort, typename	->	Bool	(* element is in a set	*)
	Union:	typename, typename	->	typename	(* union of 2 sets	*)
	Ints_:	typename, typename		typename	(* intersection of 2 sets	
	Minus:	typename, typename		typename	(* substraction of 2 sets	
	eq:	typename, typename	->	Bool	(* equality	*)
	oq: _ne_:	typename, typename		Bool	(equally	,
	lt:	typename, typename		Bool	(* lexicographical order	r *)
	Includes:	typename, typename		Bool	(* set1 \supseteq set2	*)
	IsSubsetOf	typename, typename	->	Bool	(* set1 \subseteq set2	*)
	Card:	typename	->	Nat	(* # of elements in a set	t*)
	constant1:		->	typename	(* constants for set	*)
	constant2:			typename	(* values	*)
				* 1		,
	•••					
	constantn:		->	typename		
	Auxiliary Ope	erations:				
	Insert_1:	sort, typename	->	typename	(* Implements a set as a	a *)
					(* sorted sequence	*)

ARRAY

An array is a finite sequence of fixed length with elements over the same type. The single elements of an array are accessed by an index value, specifying the number of the desired element in the array. The number of elements of an array is specified by the range of the index values. Arrays can be nested. The depth of the nesting specifies the dimension of the array, which is also called an *n*-dimensional array.

Array types are used in the specification for mathematical applications, like matrixes, as coordinates of geometrical data elements, for the encoding of characters, for the control information of low level devices, etc.

The operations that can be performed on arrays are the selection of array elements and the setting of the values of array elements. The conceptual framework provides equality and order relations on arrays.

The index range '[n..m]' has a lower limit 'n' and an upper limit 'm'. The number of array elements is 'm-n+1'. The lower border must always be less or equal to the upper one. The n-dimensional array is specified by n index ranges ('n..m').

Array Scheme	arrowtypo	typonomo				
Array Scheme	arraytype is	typename import_type [n m,, q r	1			
	elements	sort	1			
	values	constant1 = $ \langle$ expressioni,,	, exp	ressionj \rangle ;	(* OPTIONAL	*)
		$constant2 = \langle expressionk,$			(* OPTIONAL	*)
		$constantn = \langle expressionm,$., ex	pressionn \rangle ;	(* OPTIONAL	*)
	endtype					
	term and es spond with the range de ple: an array	' could be any term; it could be pecially also an array term. The its definition, i.e. the nesting and efinition. The last subrange spec y with the nesting level one has this value: $ \langle \langle 2, 4, 6, 1 \rangle , \langle 5, 2, \rangle $	only d the ifies a rar	v rule is that the number of elem the range of the age [1 2, 1 4	array constant must corr nents must be consistent we innermost array. An exa	re- with am-
Types	typename					
	to the nestir The types 't set operation	ypes: typename{_sub}*_sub, ty ng level ypename{_sub}*_sub' define th ns and equality predicates. The t d set operations with more than	ne ari Types	rays of the respe 'typename{_su	ective nesting with select	and
Library Types	DecimalNat	turalNumber				
Sorts	typename					
	Auxiliary S	orts: typename{_sub} ⁺				
Operations		pename_sub_sub [*] ,, typename rt,, sort			me_sub [*] (* constructor o*(* for the depest nestin	*) Ig *)
	nth: typ	bename, Nat,, Nat	->	sort	(* element selector	*)
	• •	bename, Nat,, Nat, sort		typename	(* set value of array-el-	
	eq: typ	bename, typename	->	Bool	(* equality	*)
	• •	bename, typename	->	Bool		
	lt: typ	bename, typename	->	Bool	(* lexicographical orde	er*)
	constant1:		->	typename	(* constants for array	*)
	constant2:		->	typename	(* values	*)
	 constantn:		->	typename		
	Auxiliary O	perations				
	nth: typ setn: typ _eq_: typ _ne_: typ _lt_: typ nth: typ	ename_sub ⁺ , Nat ename_sub ⁺ , Nat, typename_sub_sub ⁺ ename_sub ⁺ , typename_sub ⁺ ename_sub ⁺ , typename_sub ⁺ ename_sub ⁺ , Vat,, Nat ename_sub ⁺ , Nat,, Nat, sort	-> -> -> -> -> -> -> -> -> -> -> -> -> -	typename_sub_su typename_sub ⁺ Bool Bool sort typename_sub ⁺	ıb ⁺	

Signatures of Translated Structured Data Types

RECORD

A records is a collection of a fixed number of components that may be written down in any order. Each record component has its own type. The components of a record are accessed by field selectors that are associated to them. Records can have optional components, this means that these components might be omitted in the definition of a record. A field can have a default value, which is used in case a component is omitted in the definition.

Records are used in a wide range of applications. They are useful in all applications where multiple informations of different types are associated to one element. In OSI specifications, the records are used for service primitives. An other example of records are files in a data base.

Operations on records are the selection of record components and the setting of the values of record components. It is possible to test whether a component of a record is set. This test is useful before the selection of a component, because the operation is only defined if the record component is set. Another available predicate gives information, about the completeness of a record, i.e. tests whether all mandatory components are set. The conceptual framework provides equality and order relations on records.

Record Scheme	recordtype	• 1				
	is	<pre>import_typen, import_</pre>	type2	2,, import_	_type1	
	fields	selector1: sort1,				
		selector2: sort2 opt	iona	l,		
		, selectorn: sortn def a	ault	= typename	{ selectorg(expressiong),, selectorh(expressionh)}	
	values			• /	(* OPTIONAL	*)
					<pre>sioni),, selectorj(expressionj)} ; sionk),, selectorl(expressionl)} </pre>	
		constantn = typename	{sele	ectorm(expre	ssionm),, selectorn(expressionn)} :
	endtype	·····	(~	·····	,,,	·/ J •
Types	typename					
		•• •			ypename00, typename0,	
Library Types	Boolean, Nat	uralNumber, Set				
Sorts	typename					
Operations	init_typenam	e:	->	typename	(* initializes a record with the (* default setting	*) *)
	complete:	typename	->	Bool	(* tests whether all non optional (* fields are set	*) *)
	selector1:	typename	->	sort1	(* selects the value of a record	*)
	selector2:	typename	->	sort2	(* field	*)
	selectorn:	typename	->	sortn		
	set selector1	: typename, sort1		typename	(* sets the value of a field	*)
		: typename, sort2	-> ->	• •	(sets the value of a field)
		. typename, sort2		typenume		
		: typename, sortn	->	typename		
	eq:	typename, typename	->	Bool	(* equality*)	
	ne:	typename, typename				
		typename, typename			(* lexicographical order*)	
						*
	constant1:			typename	(* constants forrecord	*)
	constant2:		->	typename	(* values	*)
	constantn:		->	typename		

enumtype	typename_se	el					
is	{ selector1.	{ selector1, selector2,, selectorn }					
endtype	1	501000012	,, second ,				
uniontype	typename_co	omponen	t				
is	Boolean, typename_sel, import_typen, import_type2, import_type1						
	{ selector		sort1,				
	selector	2:	sort2,				
	, selectori	n•	sortn }				
endtype	Selector		sorur }				
enacype							
settype	typename00						
is	import_type						
elements	sort						
endtype							
type	typename0						
is	typename00						
renamedby	51						
sortnames	typename for	r typenai	me 00				
endtype							
tvpe	typename se	et interfa	ce				
type is	typename_se typename0	et_interfa	ce				
	• •		ce ne_component	->	typename_sel		
is	typename0	typenan			• •		
is	typename0 sel: add: remove:	typenan typenan typenan	ne_component ne_component, typename ne_sel, typename	e -> ->	typename typename		
is	typename0 sel: add: remove: _overlay_:	typenan typenan typenan typenan	ne_component ne_component, typename ne_sel, typename ne, typename	e -> -> ->	typename typename typename		
is	typename0 sel: add: remove:	typenan typenan typenan typenan typenan	ne_component ne_component, typename ne_sel, typename	e -> -> -> ->	typename typename		

Signatures of Translated Structured Data Types

Dependencies of Generated Types

Formal Syntax Definition of Structured Data Types

		definition shall be included into	cturd data type specifications is defined. The syntax the syntax definition of the LOTOS standard lata-type-definition' and 'term-expression' are rede-
DATA TYPE	A.1	data-type-definition =	type-symbol type-identifier is-symbol p-expression end-type-symbol structured-data-type-definition library-declaration.
STRUCTURED DATA TYPE	A.2	structured-data-type-definition =	 enumeration-type-definition union-type-definition sequence-type-definition set-type-definition array-type-definition record-type-definition.
	A.3	enumeration-type-definition =	enumeration-type-symbol sort-and-type-identifier is-symbol open-set-symbol [operation-identifier-list] close-set-symbol [subclass-expression] end-type-symbol.
	A.4	union-type-definition =	union-type-symbol sort-and-type-identifier is-symbol type-union open-set-symbol union-projection-list close-set-symbol [subclass-expression] end-type-symbol.

	A.5	sequence-type-definition =	sequence-type-symbol sort-and-type-identifier is-symbol type-identifier elements-symbol sort-identifier [initial-setting-expression] end-type-symbol.
	A.6	set-type-definition =	set-type-symbol sort-and-type-identifier is-symbol type-identifier elements-symbol sort-identifier [initial-setting-expression] end-type-symbol.
	A.7	array-type-definition =	array-type-symbol sort-and-type-identifier is-symbol type-identifier elements-symbol sort-identifier open-bracket-symbol index-list close-bracket-symbol [initial-settings-expression] end-type-symbol.
	A.8	record-type-definition =	record-type-symbol sort-and-type-identifier is-symbol type-union fields-symbol [record-projection-list] [initial-settings-expression] end-type-symbol.
SUBCLASS	A.9	subclass-expression =	subclass-symbol subclass-list.
	A.10	subclass-list =	subclass [subclass-list].
	A.11	subclass =	operation-identifier open-set-symbol [operation-identifier-list] close-set-symbol.
ARRAY INDEX	A.12	index-list =	number range-symbol number [comma-symbol index-list].
	A.13	number =	"1" "2" "3" " 4" "4" "5" "6" "7" "8" "9" { digit }.
INITIAL SETTING	A.14	initial-settings-expression =	values-symbol value-definition-list.
	A.15	value-definition-list =	value-definition [value-definition-list].
	A.16	value-definition =	operation-identifier equal-symbol value-expression.
RECORD PROJECTION	A.17	record-projection-list =	projection [feature-expression] [comma-symbol record-projection-list].
UNION PROJECTION	A.18	union-projection-list =	projection [comma-symbol union-projection-list].
	A.19	projection =	operation-identifier colon-symbol sort-identifier.
	A.20	feature-expression =	default-symbol equal-symbol value-expression optional-symbol.
AUXILIARIES	A.21	sort-and-type-identifier =	identifier.
		operation-identifier-list =	operation-identifier [comma-symbol operation-identifier-list].

SHORT VALUE EXPRESSIONS		term-expression =	 value-identifier operation-identifier [value-expression-list] open-parenthesis-symbol value-expression close-parenthesis-symbol set-expression sequence-expression array-expression record-expression.
	A.24	set-expression =	open-set-symbol [value-expression-list] close-set-symbol.
	A.25	sequence-expression =	open-sequence-symbol [value-expression-list] close-sequence-symbol.
	A.26	array-expression =	open-array-expression value-expression-list close-array-symbol.
	A.27	record-expression =	operation-identifier open-set-symbol record-assignment-list close-set-symbol.
	A.28	record-assignment-list =	operation-identifier open-parenthesis-symbol term-expression close-parenthesis-symbol [comma-symbol record-assignment-list].
WORD SYMBOLS	A.29	enumeration-type-symbol =	"enumtype".
	A.30	union-type-symbol =	"uniontype".
	A.31	sequence-type-symbol =	"seqtype".
	A.32	set-type-definition =	"settype".
	A.33	array-type-symbol =	"arraytype".
	A.34	record-type-symbol =	"recordtype".
	A.35	subclass-symbol =	"subclass".
	A.36	fields-symbol =	"fields".
	A.37	optional-symbol =	"optional".
	A.38	default-symbol =	"default".
SPECIAL SYMBOLS	A.39	open-set-symbol =	" {".
	A.40	close-set-symbol =	" }".
		open-sequence-symbol =	" (".
		close-sequence-symbol =	")]".
		open-array-symbol =	" <".
		close-array-symbol =	"> ".
		range-symbol =	"
	A.46	reverse-arrow-symbol =	"<-".

Installation of the Structured Data Type Translation Tool 'pao2ao'

FILES	'pao2ao' stands for 'powerfull ACT ONE to ACT ONE'. The following files contain the source code for the pao2ao tool:		
	makefile trans_main.c trans_main.h trans_1.1 trans_y.y trans_decl.k trans_array.k trans_lib.k trans_record.k trans_rw.k trans_unp.k trans_graph.k	makefile to create translator main function for translator type declarations of variables and functions lex file (token definitionen) yacc file (syntax definition) kimwitu tree declarations rewrite functions for array types and expressions functions to create the library types rewrite functions for record types and expressions general rewrite function pool unparse functions to create text file of the kimwitu tree generation of graphical annnotations for visualization with DEMON	
MAKEFILE The makefile shows which variables have to build:		which variables have to be defined and how the executable will be	
	IT = KFILES = CC = YOURFILES = ALLOBJS = GENERATED_C = GENERATED_H =	trans_ \${IT}decl.k \${IT}rw.k \${IT}lib.k \${IT}unp.k \${IT}record.k\ \${IT}array.k \${IT}graph.k cc \${KFILES} \${IT}y.y \${IT}l.l \${IT}main.c k.o rk.o csgiok.o unpk.o\ \${KFILES:k=o} \${IT}y.o \${IT}l.o \${IT}main.o k.c rk.c csgiok.c unpk.c \${KFILES:k=c} k.h rk.h csgiok.h unpk.h \${KFILES:k=h}	

GENERATED_BY_KC YACC = YFLAGS = KC = GENERATED_LN =	= \${GENERATED_C} \${GENERATED_H} bison -dyv /net/jupiter/usr85/SEM/kimwitu/kc-distr.V3_8/src/Gen/kc \${IT}decl.ln \${IT}rw.ln \${IT}lib.ln \${IT}unp.ln\ \${IT}record.ln \${IT}array.ln \${IT}graph.ln\ k.ln rk.ln csgiok.ln unpk.ln \${IT}main.ln \${IT}y.ln \${IT}1.ln
new_\${IT}:	\${ALLOBJS} \${CC} \${CFLAGS} \${ALLOBJS} -ll -0 \$@
\${GENERATED_BY_K	
kctimestamp:	\${KFILES}
	<pre>\${KC} \${KFILES}; touch kctimestamp</pre>
\${ALLOBJS}:	k.h
\${IT}main.o \${IT}l.o:	x.tab.h
\${IT}main.o \${KFILES:	
\${IT}main.o rk.o:	rk.h
\${IT}main.o csgiok.o:	csgiok.h
\${IT}main.o unpk.o:	unpk.h
x.tab.h:	y.tab.h
	-cmp -s x.tab.h y.tab.h cp y.tab.h x.tab.h
lint:	\${GENERATED_LN}
	-@ lint -u -n -q -v \${CFLAGS} \${GENERATED_LN} 2>&1
	sed -e '/warning:/d' -e '/malloc[,:]/d' -e '/printf[,:]/d'\ -e '/scanf[,:]/d' -e '/^free[,:]/d'
.c.ln:	
.0.111.	lint -u -n -q -v -i \$< 2>&1 sed '/warning:/d'
	$\lim_{n \to \infty} u = \frac{1}{\sqrt{2}} \sqrt{2} \sqrt{2} \sqrt{2} \sqrt{2} \sqrt{2} \sqrt{2} $

Annex Signatures of Library Types

	The following signatures are provided for those that do not use provides the sorts, operations and equations of the three librar				
	•	his library is the IS 8807 Annex A standard libray, modified by LotosPhere WP2, and commended for use with tools. the changes are in types Set and SetElement. bb 1991			
NaturalNumber	type NaturalNumber is Boolean sorts Nat opns 0 Succ $_+_,_*_,_**_$ $_eq_,_ne_,_lt_,_le_,_ge_,_$ eqns forall m, n : Nat ofsort Nat m + 0 = m; m + Succ(n) = Succ(m) + n; m * 0 = 0;	: -> Nat : Nat -> Nat : Nat, Nat -> Nat gt_ : Nat, Nat -> Bool			
	m * Succ(n) = m + (m * n); $m ** 0 = Succ(0);$ $m ** Succ(n) = m * (m ** n);$ ofsort Bool $0 eq 0 = true;$ Succ(m) eq 0 = false; m ne n = not(m eq n); 0 lt 0 = false; Succ(n) lt 0 = false; m le n = m lt n or (m eq n); m ge n = not(m lt n); m gt n = not(m le n); endtype	0 eq Succ(m) = false; Succ(m) eq Succ(n) = m eq n; 0 lt Succ(n) = true; Succ(m) lt Succ(n) = m lt n;			

String

tuno	String
type is	Boolean, Element, NaturalNumber
	String
sorts	6
opns	String
	+
	+
	++
	Reverse
	Length
	\Leftrightarrow
	eq, _ne_
eqns	forall s, t : String, x, y : Element
ofsort	String
	String(x) + y = x + String(y);
	x + s + y = x + (s + y);
	String(x) ++ s = x + s;
	x + s ++ t = x + (s ++ t);
	Reverse(String(x)) = String(x);
	Reverse(x + s) = Reverse(s) + x;
	String(x) = x + <>;
	<>+ x = x + <>;
	<> ++ s = s;
	(* new equation *)
	s ++ <> = s;
	(* end new equation *)
	Reverse(<>) = <>;
ofsort	Nat
	Length(String(x)) = Succ(0);
	Length(x + s) = Succ(Length(s));
	$\text{Length}(\ll) = 0;$
ofsort	Bool
	<> eq <> = true;
	$\langle \rangle$ eq (x + s) = false;
	$x + s eq \ll = false;$
	x eq y => x + s eq (y + t) = s eq t;
	x ne y $=> x + s$ eq (y + t) = false;
	s ne $t = not(s eq t);$
endtype	· • •

: Element	-> String
: Element, String	-> String
: String, Element	-> String
: String, String	-> String
: String	-> String
: String	-> Nat
:	-> String
: String, String	-> Bool

Set

type is	Set SetElement, Boolean, NaturalNumber				
sorts	Set			G .	
opns	{} }			-> Set	
	Insert, Remove, Insert_1		: Element, Set	-> Set	
	IsIn, _NotIn_		: Element, Set	-> Bool	
	Union, _Ints_, _Minus_		: Set, Set	-> Set	
	eq, _ne_, _lt_, _Includes_, _IsSubsetOf_ Card		: Set, Set : Set	-> Bool -> Nat	
eqns	forall x, y : Element, s, t : Set				
ofsort	Set $I_{\text{mont}}(y, \beta) = I_{\text{mont}}(y, \beta)$				
	Insert(x, $\{\}$) = Insert_1(x, $\{\}$); v by $=$ Insert(x, Insert 1(x, z)) = Insert 1(x)	Incont 1	(m. a)).		
	x lt y => Insert(x, Insert_1(y, s)) = Insert_1(x, s)	, insert_1	y, s));		
	$Insert(x, Insert_1(x, s)) = Insert_1(x, s);$	T 44	~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~		
	y lt x => Insert(x, Insert_1(y, s)) = Insert_1(y	, Insert(x,	s));		
	Remove $(x, \{\}) = \{\};$				
	Remove(x, Insert_1(x, s)) = s;				
	$x \text{ lt } y \Rightarrow \text{Remove}(x, \text{Insert}_1(y, s)) = \text{Insert}_1(y, s)$				
	y lt x => Remove(x, Insert_1(y, s)) = Insert_	1(y, Remo	ve(x, s));		
	$\{\}$ Union s = s;				
	Insert_1(x, s) Union $\{\}$ = Insert_1(x, s);				
	x lt y => Insert_1(x, s) Union Insert_1(y, t) =				
	y lt x => Insert_1(x, s) Union Insert_1(y, t) = Insert_1(y, Insert_1(x, s) Uni on t);				
	Insert_1(x, s) Union Insert_1(x, t) = Insert_1	(x, s Unio	n t);		
	$\{\}$ Ints s = $\{\};$				
	Insert_1(x, s) Ints $\{\} = \{\};$				
	Insert_1(x, s) Ints Insert_1(x, t) = Insert_1(x, t)	, s Ints t);			
	x lt y => Insert_1(x, s) Ints Insert_1(y, t) = s $\frac{1}{2}$	Ints Insert	1(y, t);		
	y lt x => Insert_1(x, s) Ints Insert_1(y, t) = In				
	s Minus { } = s;				
	s Minus Insert_ $1(x, t) = \text{Remove}(x, s)$ Minus	t;			
ofsort	Bool				
	$x \text{ IsIn } \{ \} = \text{false};$				
	x IsIn Insert_1(x, s) = true;				
	y lt x => x IsIn Insert_ $1(y, s) = x$ IsIn s;				
	x lt y => x IsIn Insert(y, s) = false;				
	x NotIn $s = not(x IsIn s);$				
	s Includes { } = true;				
	s Includes Insert_ $1(x, t) = x$ IsIn s and (s Incl	udes t).			
	s IsSubsetOf $t = t$ Includes s;	luues t),			
	{} eq {} = true;				
	$\{\}$ eq Insert_1(x, s) = false;				
	Insert_1(x, s) = false; f = false;				
	$x eq y => Insert_1(x, s) eq Insert_1(y, t) = s et x eq y => Insert_1(y, s) eq Insert_1(y, t) = for$				
	x ne y => Insert_1(x, s) eq Insert_1(y, t) = fat	ise,			
	$\{\}$ It $\{\}$ = false;				
	{} It Insert_1(x, s) = true;				
	Insert_1(x, s) It $\} = $ false;				
	x lt y => Insert_1(x, s) lt Insert_1(y, t) = true	;			
	Insert_1(x, s) lt Insert_1(x, t) = s lt t;				
	y lt x => Insert_1(x, s) lt Insert_1(y, t) = false	e;			
c	s ne $t = not(s eq t);$				
ofsort	Nat				
	$\operatorname{Card}(\{\}) = 0;$				
	$Card(Insert_1(x, s)) = Succ(Card(s));$				
endtype					

Auxiliary Types	type is	Boolean		
	sorts	Bool		
	opns	true, false	:	-> Bool
		not	: Bool	-> Bool
		and, _or_, _xor_, _implies_, _iff_, _eq_, _ne_	: Bool, Bool	-> Bool
	eqns	forall x, y : Bool	,	
	ofsort	Bool		
		not(true) = false;		
		not(false) = true;		
		x and true = x;		
		x and false = false;		
		x or true = true;		
		x or false = x;		
		x xor $y = x$ and not(y) or (y and not(x));		
		x implies $y = y$ or $not(x)$;		
		x iff $y = x$ implies y and (y implies x);		
		x eq y = x iff y;		
	x ne y = x xor y;			
	endtype			
	type	FBoolean		
	is			
	formalso	rts FBool		
	formalop	ons true	:	-> FBool
		not	: FBool	-> FBool
	formaleq			
	ofsort	FBool		
		not(not(x)) = x;		
	endtype			
	type	Element		
	is	FBoolean		
	formalso			
	formalop		: Element, Element	-> FBool
	formaleq			
	ofsort	Element		
	C , ,	$x eq y \Rightarrow x = y;$		
	ofsort	FBool		
		$x = y \Rightarrow x eq y = true;$		
	an dtrin a	x ne $y = not(x eq y);$		
	endtype			
	type SetElement is Element			
	formalop	ons _lt_	: Element, Element	nt -> FBool
	endtype			