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Generalising Feature Interactions in Email

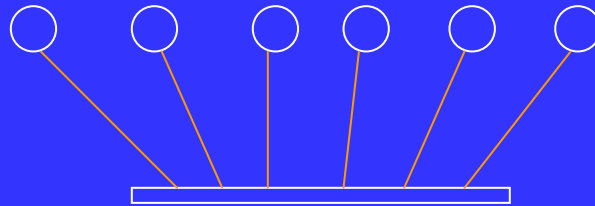
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Motivation

- property based approach to feature interaction analysis
- interaction analysis should
 - *be automated*
 - *generalise to systems containing any number of components*
- based on Hall's email model from FIW'00

Email system

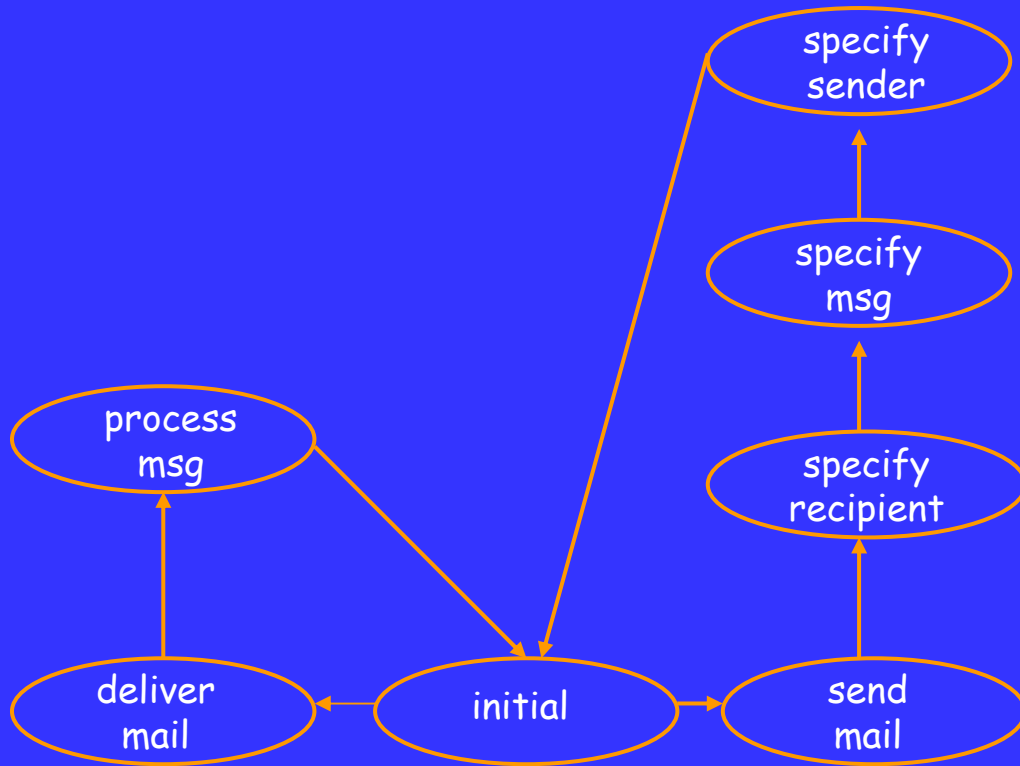
Client server architecture



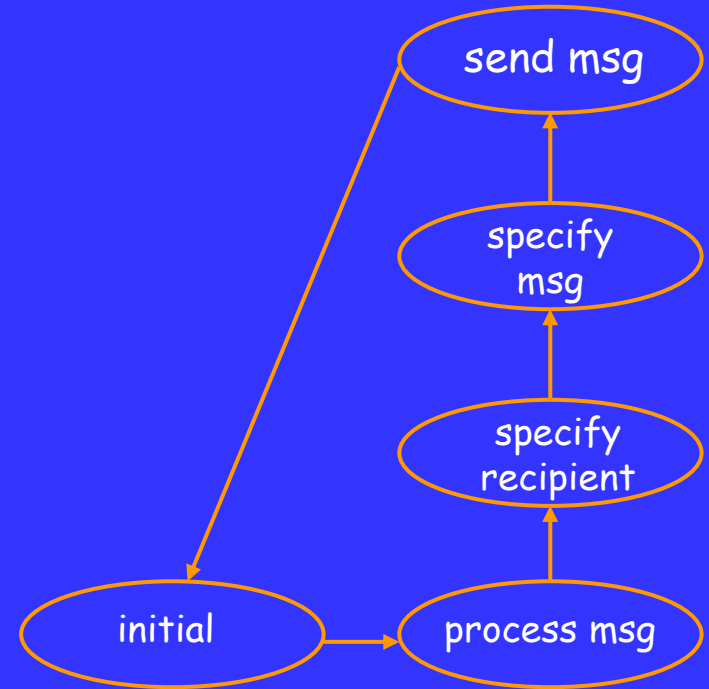
basic email + features at client *or* server

- encrypt -- key of intended recipient
- decrypt - key of actual recipient
- filter - msgs from address
- forward -msgs to address
- autorespond -to first incoming msgs

Email system



client process



mailer process

Promela implementation

Clients and the Mailer are processes.

- communication is asynchronous
- channels associated with each client and the mailer
- delivery of mail takes precedence over sending
- busy waiting
- tension between
 - atomicity/number of variables/level of abstraction

Property based approach

Example Properties in linear temporal logic:

- *messages are delivered to intended recipients*

$\square(p \parallel q)$ where $p = (\text{last_del_to}_i_to = i)$
 $q = (\text{last_del_to}_i_to = M)$

- *messages are forwarded, client_i has forwarding to client_j*

$\exists \leftrightarrow (\neg(p \parallel q))$ where $p = (\text{last_del_to}_j_to = M)$
 $q = (\text{last_del_to}_j_to = j)$

(i.e. client_j can receive mail not addressed to j)

observation variables: $\text{last_del_to}_i_to, \text{last_del_to}_i_body, \text{last_sent_from}_i_to$

Property based approach

Feature interaction analysis based on:

$$f_0 \parallel \text{System} \models \phi \quad \text{but} \quad f_0 \parallel f_1 \parallel \text{System} \not\models \phi$$

for example:

$$\text{Client}_0 \parallel \text{Client}_1 \parallel \text{Client}_2 \parallel \text{Client}_3 \parallel \text{Mailer} \models \phi$$

but

$$\text{Client}'_0 \parallel \text{Client}_1 \parallel \text{Client}_2 \parallel \text{Client}_3 \parallel \text{Mailer} \not\models \phi$$

Reasoning

- Use model-checker (SPIN) for reasoning
- Results take the form

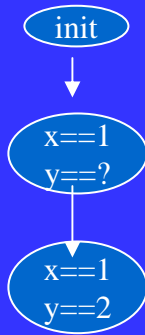
$$\mathcal{M}(p_0 \parallel p_1 \parallel p_2 \parallel p_3 \parallel \text{mailer}) \models \phi(0,1,\dots,t)$$

where

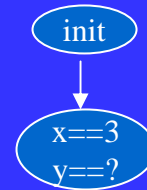
- $p_0 \dots p_3$ are instances of a *parameterised* process p
- $p_0 \dots p_3$ are not, in general, isomorphic
- ϕ is temporal logic formula containing free variables indexed by $0,1,\dots,t$ $0 \leq t < 3$
- $\mathcal{M}(\dots)$ is the model (Kripke structure) of the concurrent processes

model checking

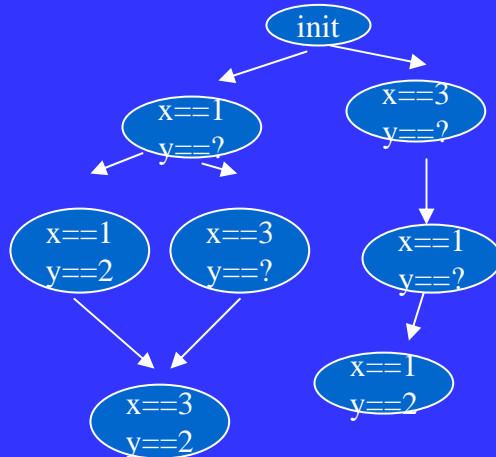
proctype A()
{x=1;y=2}



proctype B()
{x=3}



program:
byte x,y;
init {run A(); run B();}

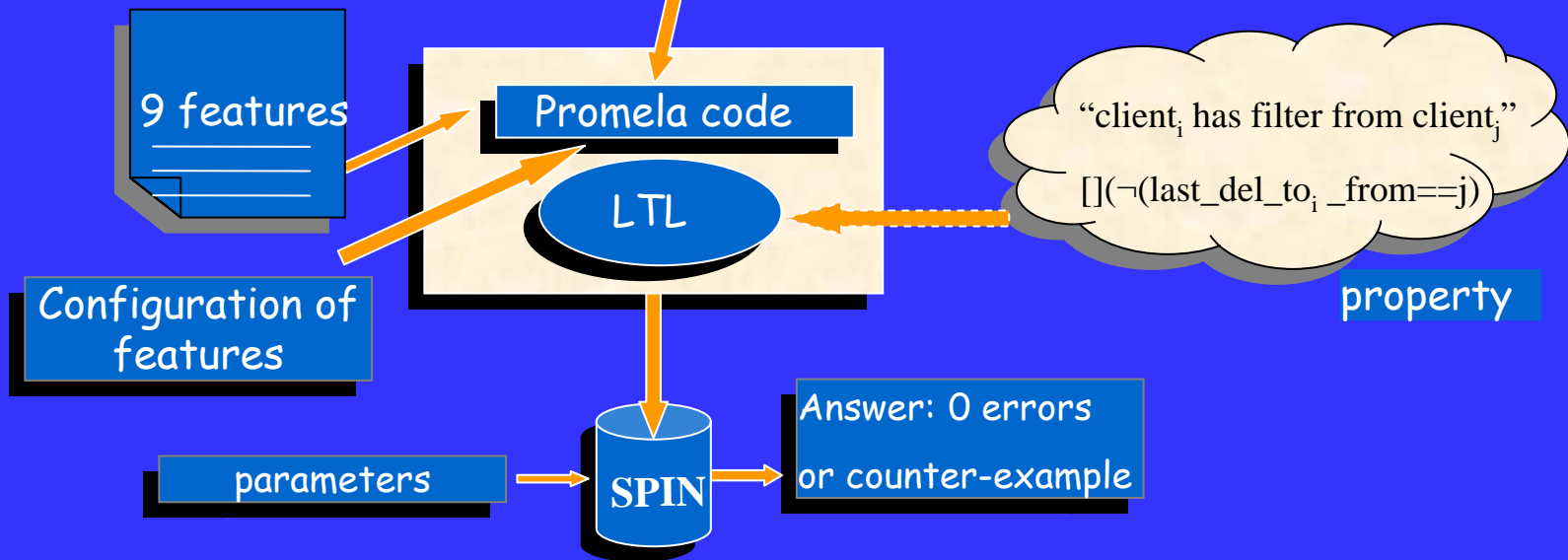
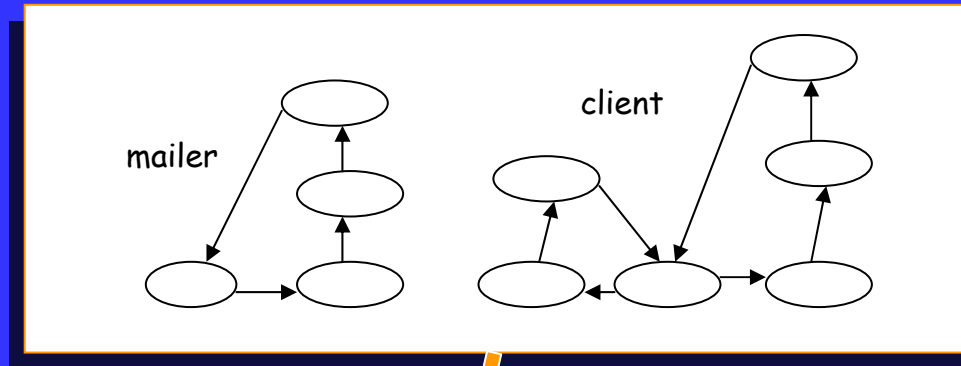


A program/system is the asynchronous interleaving product of the automata.

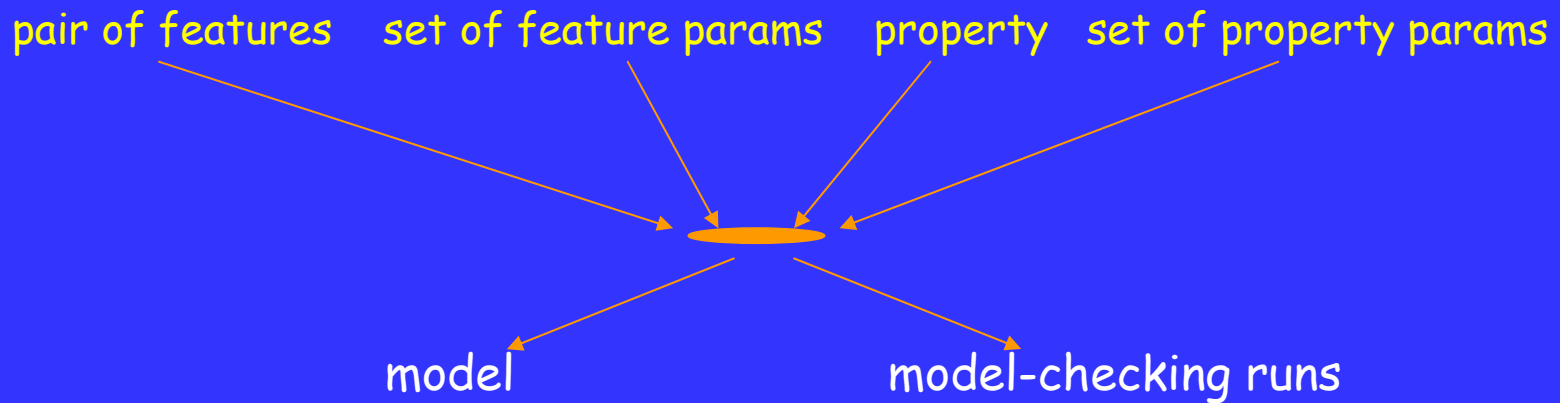
Reason about satisfaction of LTL formula by considering synchronous product of system with never claim.

overall approach

Mailer and client basic behaviour



automation



up to 5 client processes required for this feature set
111 feasible parameter sets, after symmetry reduction

via perl scripts

interactions

- the usual!
- both single user and multiple user
- multiple user agree with Hall results

The interesting question is.. do results scale?

- do they hold for 7 processes, 8 processes, ... ?
- for every problem, we eventually run out of memory (or patience)
 - usually around 6 processes

Generalisation

- What we really want is to show is

$$\forall n. \mathcal{M}(p_0 \parallel p_1 \parallel p_2 \parallel \dots \parallel p_{n-1} \parallel \text{mailer}) \models \phi(0,1,\dots,t)$$

Not possible with model checking

Undecidable (Apt + Kozen, 1986)

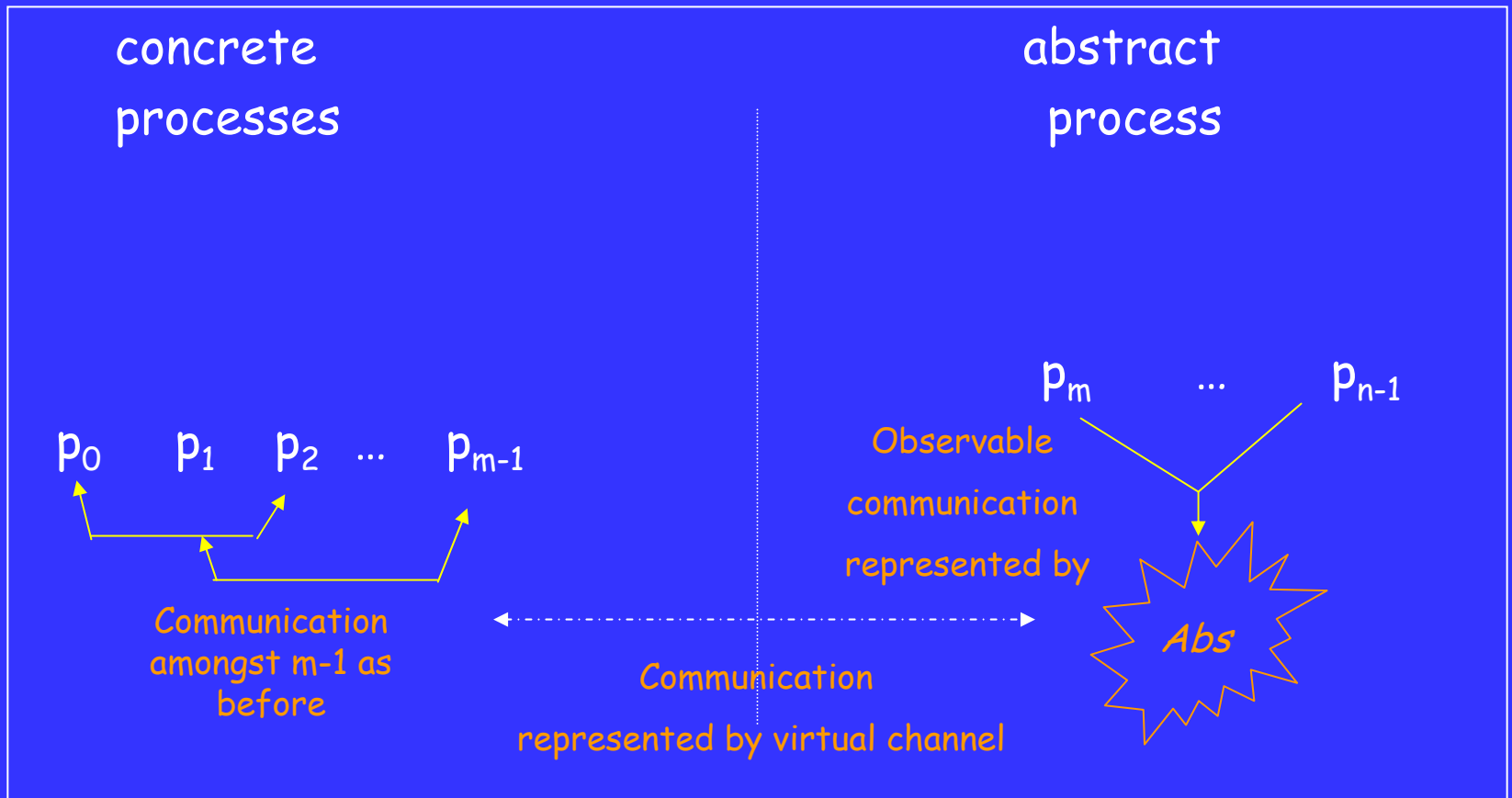
- Induction is very hard
- Is abstraction possible?

Abstraction

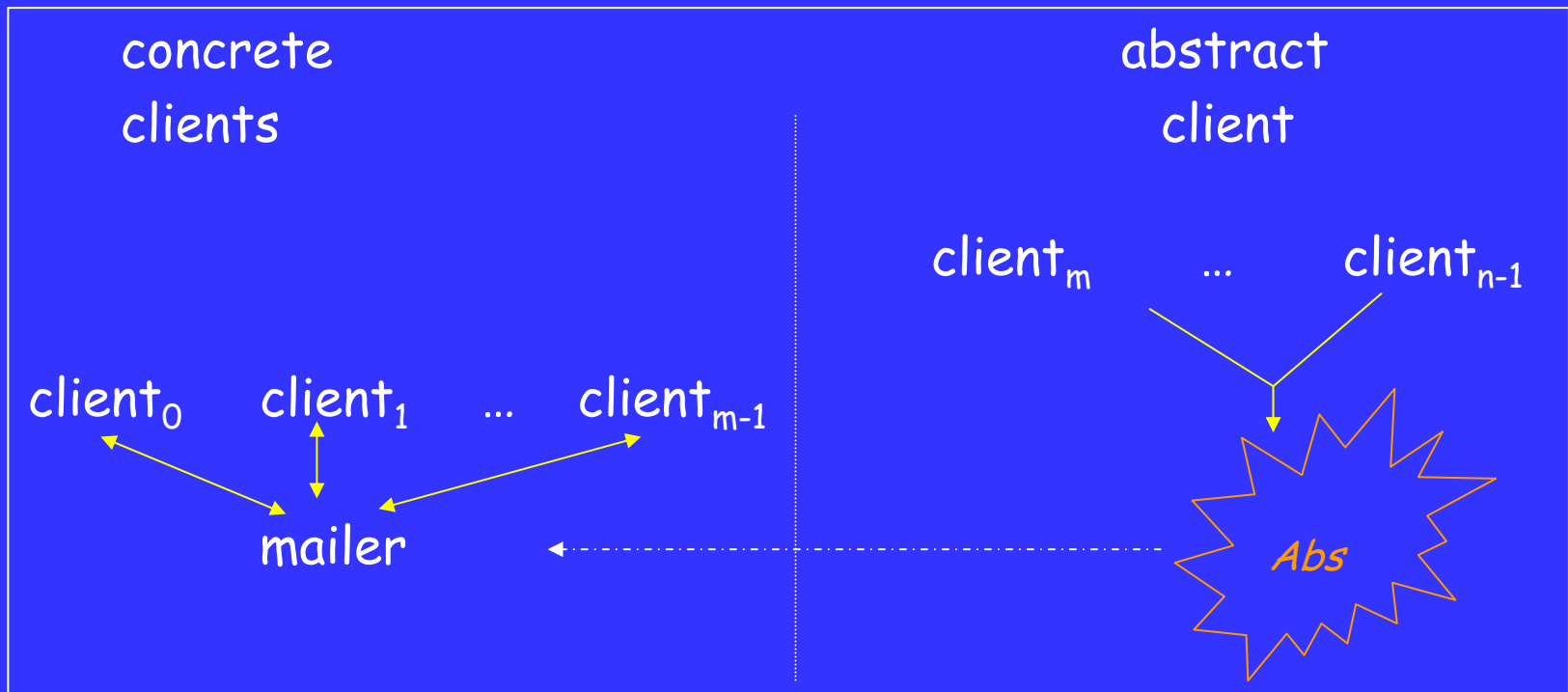
- What do we mean by abstraction?
 - not an abstraction of one system, but of a family of systems
 - Choose appropriate constant m such that $t \leq m-1$
 - $p_0, p_1 \dots p_{m-1}$ are *concrete*
 - $p_m, p_{m+1} \dots p_{n-1}$ are *abstract*
 - Represent the most *general, observable* behaviour of the abstract processes $p_m \parallel p_{m+1} \parallel \dots \parallel p_{n-1}$ by a process *Abs*
 - Modify the interaction between concrete and abstract processes, i.e. modify the concrete processes.

All processes should be generated automatically from p .

Abstract approach



Abstract approach: email



- Only "read" behaviour from *Abs*. (Choice rather than actual read)
- mailer "writes" to *Abs* if not blocked. (Choice).
- Not strictly a conservative extension. *Abs* can send mail if there is mail to to be delivered. Does not affect functional behaviour.

Results

Checking done with perl scripts.

No new interactions (not surprising, given feature set).

Complexity lies between that of m and $m+1$ (concrete) clients.

m depends on feature parameters and property parameters (essentially union).

Soundness

The $p_m, p_{m+1}, \dots, p_{n-1}$ need not be isomorphic, or even observationally equivalent, but we make some assumptions about both concrete and abstract processes:

1. All interaction is through communication channels.
2. All processes are *open symmetric* - behave the same with respect to isomorphic processes - no integer literals or constants in boolean conditions
 - $g?x; x==9 \Rightarrow \text{goto label}$ **NO**
 - $g?x; x==\text{var}_i \Rightarrow \text{goto label}$ **YES**

Theorem

$M(\text{client}_0 \parallel \text{client}_1 \parallel \dots \parallel \text{client}_{m-1} \parallel \text{mailer}' \parallel \text{Abs}) \models \phi(0,1,\dots,t)$

\Rightarrow

$\forall n. M(\text{client}_0 \parallel \text{client}_1 \parallel \dots \parallel p_{n-1} \parallel \text{mailer}) \models \phi(0,1,\dots,t).$

Proof

Show *simulation*. Depends on way we construct *Abs* and *mailer'*: consider how to match

- concrete process reads from abstract channel
- concrete process write to abstract channel

...

- abstract process reads from concrete channel
- abstract process writes to abstract channel

Conclusions

- property based approach to interaction analysis
 - automated using perl scripts - to tailor model and generate runs.
- abstraction to generalise results about infinite *families* of communicating processes
 - processes are not isomorphic.
- generation of abstract model is straightforward
 - implemented in perl scripts
 - lower bound for m
 - for this feature set abstraction approach is tractable.

further work - is the *abstraction approach* constructing an *invariant*?