Design of Secure Computer Systems CSI4138/CEG4394 Notes on the Modular Arithmetics and Galois Fields

1 Congruence and modular arithmetics

Let a, b, and n be non-negative integers, i.e. $n \in \mathcal{N}$ the set of natural numbers, and $n \neq 0$; then a is said to be congruent to b modulo n, that is

$$a \equiv_n b$$
 if and only if, $a - b = kn$

for some integer k. In other words, n divides the difference (a - b). For instance,

$$17 \equiv_5 7$$
 since $17 - 7 = 2 \times 5$.

b is a residue of a modulo n and also a is a residue of b modulo n. For any modulus n, the set of integers $\{0, 1, \ldots, n-1\}$ forms a complete set of residues modulo n:

$$\{r_1, \ldots, r_n\} = \{0, 1, \ldots, n-1\}$$

The residue r of a modulo n is in the range [0, n-1]. Note that

 $a \mod n = r \qquad \Rightarrow \qquad a \equiv_n r \qquad \text{but not the converse:}$

 $a \equiv_n r \quad \not\Rightarrow \quad a \mod n = r$

meaning that $a \equiv_n r$ does not imply that $a \mod n = r$; for instance,

 $17 \mod 5 = 2 \quad \Rightarrow \quad 17 \equiv_5 2 \quad \text{but}$

$$17 \equiv_5 7 \quad \Rightarrow \quad 17 \mod 5 = 7$$

1.1 Properties of modular arithmetics:

Let the symbol (\odot) represent either an addition (+) or a multiplication (\times) operation.

1. Existence of identities:

 $a + 0 \mod n = 0 + a \mod n = a$ $a \times 1 \mod n = 1 \times a \mod n = a$ 2. Existence of inverses:

$$a + (-a) \mod n = 0$$

 $a \times (a^{-1}) \mod n = 1$ if $a \neq 0$

3. Commutativity:

$$a \odot b \mod n = b \odot a \mod n$$

4. Associativity:

 $a \odot (b \odot c) \mod n = (a \odot b) \odot c \mod n$

5. Distributivity:

$$a \times (b+c) \mod n = [(a \times b) + (a \times c)] \mod n$$

6. Reducibility:

 $(a \odot b) \mod n = [(a \mod n) \odot (b \mod n)] \mod n \quad \text{or equivalently:}$ $(a+b) \mod n = [(a \mod n) + (b \mod n)] \mod n$ $(a \times b) \mod n = [(a \mod n) \times (b \mod n)] \mod n$

- Ring: associativity and distributivity
- Commutative ring: associativity, distributivity, and commutativity
- Galois field: commutative ring where each element $\neq 0$ has a multiplicative inverse.

2 Principle of modular arithmetics (reducibility)

The reducibility property states that:

$$(a \odot b) \mod n = [(a \mod n) \odot (b \mod n)] \mod n$$

Proof:

Two integer numbers a_1 and a_2 can be written as: $a_1 = k_1n + r_1$ and $a_2 = k_2n + r_2$, where $r_1, r_2 \in [0, n-1]$, and both k_1 and k_2 are positive integers. The reducibility property can be proven for the addition operation (\odot : +) as follow:

$$(a_1 + a_2) \mod n = [(k_1n + r_1) + (k_2n + r_2)] \mod n$$

= $[(k_1 + k_2)n + r_1 + r_2)] \mod n$
= $(r_1 + r_2) \mod n$
 $(a_1 + a_2) \mod n = [(a_1 \mod n) + (a_2 \mod n)] \mod n$

by definition of a residue. Similarly, for the multiplication operation, i.e. $(\odot : \times)$:

$$(a_1 \times a_2) \mod n = [(k_1n + r_1) \times (k_2n + r_2)] \mod n$$

= $[(k_1k_2n^2) + (k_1nr_2) + (k_2nr_1) + (r_1r_2)] \mod n$
= $[(k_1k_2n + k_1r_2 + k_2r_1)n + (r_1r_2)] \mod n$
= $(r_1 \times r_2) \mod n$
 $(a_1 \times a_2) \mod n = [(a_1 \mod n) \times (a_2 \mod n)] \mod n$

Principle of modular arithmetics

 $(a_1 \mod n), (a_2 \mod n)$ \longrightarrow reduction modulo $n \longrightarrow$ a_1, a_2 \odot \odot $\rightarrow \text{ reduction modulo } n \longrightarrow [(a_1 \mod n) \odot (a_2 \mod n)] \mod n$ $a_1 \odot a_2$

Modular exponentiation 3

Using the properties of modular arithmetics, modular exponentiation can be performed with the advantage of limiting the range of intermediate values:

$$e^{t} \mod n = [e \times e \times \ldots \times e] \mod n$$
$$= \{\underbrace{[e \mod n] \ [e \mod n] \ \ldots \ [e \mod n]]}_{t \text{ times}} \} \mod n$$

The intermediate values $[e \mod n]$ being reduced within the range of the modulus, that is $[e \mod n] \in [0, n-1].$

$$e^t \mod n = [\prod_{i=1}^t (e \mod n)] \mod n$$

$\begin{array}{c} \mathbf{Example}(modular\ exponentiation):\\ \text{Compute the following:}\ 11^{207}\ \mathrm{mod}\ 13 \end{array}$

$11^{207} \mod 13$	3 =	$\left[11^{128+64+8+4+2+1}\right] \bmod 13$
$11^{207} \mod 13$	3 =	$[11^{128} \times 11^{64} \times 11^8 \times 11^4 \times 11^2 \times 11] \mod 13$
$11^{207} \mod 13$	3 =	$\left\{ \left[11^{128} \bmod 13\right] \left[11^{64} \bmod 13\right] \left[11^8 \bmod 13\right] \left[11^4 \bmod 13\right] \left[11^2 \bmod 13\right] \times 11 \right\} \bmod 13$
$11^{207} \mod 13$	3 =	$\left\{ \left[11^{128} \mod 13 \right] \left[11^{64} \mod 13 \right] \left[11^8 \mod 13 \right] \left[11^4 \mod 13 \right] \times 4 \times 11 \right\} \mod 13$
$11^{207} \mod 13$	3 =	$\left\{ \left[11^{128} \mod 13 \right] \left[11^{64} \mod 13 \right] \left[11^8 \mod 13 \right] \times 3 \times 4 \times 11 \right\} \mod 13$

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\begin{array}{rcl} 11^{207} \bmod 13 & = & \left\{ \left[ 11^{128} \bmod 13 \right] \left[ 11^{64} \bmod 13 \right] \times 9 \times 3 \times 4 \times 11 \right\} \bmod 13 \\ 11^{207} \bmod 13 & = & \left\{ \left[ 11^{128} \bmod 13 \right] \times 3 \times 9 \times 3 \times 4 \times 11 \right\} \bmod 13 \\ 11^{207} \bmod 13 & = & \left\{ 9 \times 3 \times 9 \times 3 \times 4 \times 11 \right\} \bmod 13 \\ 11^{207} \bmod 13 & = & \left\{ 32076 \right\} \bmod 13 \\ 11^{207} \bmod 13 & = & 5 \end{array}
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4 Multiplicative inverses

Let $a \in [0, n-1]$ and $x \in [0, n-1]$ be a multiplicative inverse of a such that:

 $ax \mod n = 1$

a has a unique multiplicative inverse modulo n when a and n are relatively prime or, in other words, if gcd(a, n) = 1 (gcd(a, n): greatest common divisor of a and n).

Example(multiplicative inverses): Let a = 3 and n = 5, then gcd(a, n) = 1:

> $a \times i \mod 5$ $3 \times 0 \mod 5 = 0$ $3 \times 1 \mod 5 = 3$ $3 \times 2 \mod 5 = 1$ $3 \times 3 \mod 5 = 4$ $3 \times 4 \mod 5 = 2$

There is a unique inverse for each value of a. The set of inverses $\{a_i^{-1}\}$ is in fact a permutation of the set of indices $\{i\}$. Now, changing n to n = 6:

 Since $gcd(a, n) \neq 1$, the inverses of a are not unique.

If gcd(a, n) = 1, then there exists an integer x, 0 < x < n, such that:

 $ax \mod n = 1$

where, as stated above, the set $\{a \times i \mod n\}$ is a permutation of $\{i\}$. The Euclid's algorithm can be used to compute to compute the greatest common divisor of a and n.

5 Euclid's algorithm

The following algorithm determines the greatest common divisor of two numbers, e.g. a and b:

$$a = b q_1 + r_1, \quad \text{for } 0 < r_1 < b$$

$$b = r_1 q_2 + r_2, \quad \text{for } 0 < r_2 < r_1$$

$$r_1 = r_2 q_3 + r_3, \quad \text{for } 0 < r_3 < r_2$$

$$r_2 = r_3 q_4 + r_4, \quad \text{for } 0 < r_4 < r_3$$

$$\vdots$$

$$r_{k-2} = r_{k-1} q_k + r_k, \quad \text{for } 0 < r_k < r_{k-1}$$

$$r_{k-1} = r_k q_{k+1}$$

The last remainder, r_k , is the greatest common divisor of a and b, i.e. $gcd(a,b) = r_k$.

Example (gcd(a, b) using the Euclid's algorithm):

For a = 360 and b = 273, determine their greatest common divisor gcd(a, b) by employing the Euclid's algorithm.

$$360 = 273 \times 1 + 87$$

$$273 = 87 \times 3 + 12$$

$$87 = 12 \times 7 + 3$$

$$12 = 3 \times 4$$

Therefore, the greatest common divisor gcd(360, 273) is equal to the remainder $r_3 = 3$. In fact, a and b can be written as:

$$360 = 5 \times 3 \times 3 \times 2 \times 2 \times 2, \text{ and}$$

$$273 = 13 \times 7 \times 3$$

6 Inverse computation

Consider the *complete set* $\{r_i\}$ of residues modulo n:

$$\{r_1, \ldots, r_i, \ldots, r_n\} = \{0, \ldots, n-1\}$$

where r_i is a residue, such that $a \equiv_n r_i$. The *reduced set* of residues modulo n is defined as the subset of $\{r_i\}_{i=1,\dots,n}$, such that r_i is relatively prime to n (excluding 0):

$${r_i}_{i=1,...,\phi(n)}$$

where $\phi(n)$ (called Euler totient function of n) represents the number of elements in this reduced set of residues. If

$$gcd(a,n) = 1$$
 then $gcd(ar_i,n) = 1$

for the reduced set of residues $\{r_1, \ldots, r_{\phi(n)}\}$, then since (ar_i) is relatively prime with n:

$$(ar_i) \mod n = r_i$$

In other words, the set $\{r_j\}$ is a permutation of the set $\{r_i\}$:

$$\{r_j\} = \{(ar_i) \bmod n\}_{i=1,\dots,\phi(n)} = P \circ \{r_i\}_{i=1,\dots,\phi(n)}$$

The following examples give the Euler totient function $\phi(n)$ for different values of n. For instance, if n is prime then, by definition: $\phi(n) = n - 1$. For n = pq where p and q are primes:

$$\phi(n) = \phi(pq)$$

$$\phi(n) = (p-1) (q-1)$$

Examples (Euler totient function $\phi(n)$):

For the following examples, let p, q and p_i be prime numbers while e and e_i are positive integers.

1. If n = p, then the reduced set of residues is:

$$\{r_i\} = \{1, 2, \dots, p-1\}$$

whereas the Euler function is equal to:

$$\phi(n) = \phi(p) = p - 1$$

2. If $n = p^2$, the reduced set of residues is:

$$\{r_i\} = \{1, 2, \dots, p-1, p+1, \dots, 2p-1, 2p+1, \dots, p^2-1\}$$

and,

$$\phi(n) = \phi(p^2) = p(p-1)$$

3. If n = pq, the reduced set of residues is:

$$\{r_i\} = \{1, 2, \dots, pq - 1\} - \{p, 2p, \dots, (q - 1)p\} - \{q, 2q, \dots, (p - 1)q\}$$
$$\phi(n) = \phi(pq) = (pq - 1) - (q - 1) - (p - 1) = (p - 1)(q - 1)$$

4. If $n = p^e$, the reduced set of residues is:

$$\{r_i\} = \{1, 2, \dots, p^e - 1\} - \{p, 2p, \dots, (p^{e-1} - 1)p\}$$

$$\phi(n) = \phi(p^e) = (p^e - 1) - (p^{e-1} - 1) = (p^{e-1})(p - 1)$$

5. If $n = \prod_{i=1}^{t} p_i^{e_i}$, the Euler function is:

$$\phi(n) = \phi\left[\prod_{i=1}^{t} p_i^{e_i}\right] = \prod_{i=1}^{t} p_i^{(e_i-1)} (p_i-1)$$

An integer n can always be expressed as a product of primes numbers:

$$n = p_1^{e_1} \times p_2^{e_2} \times \ldots \times p_t^{e_t}$$
$$n = \prod_{i=1}^t p_i^{e_i}$$

where the p_i 's are t distinct prime numbers and their exponents e_i are positive integers. As indicated above, the number of elements in the reduced set is given by:

$$\phi(n) = \prod_{i=1}^{t} p_i^{(e_i-1)} \ (p_i-1)$$

6.1 Euler's generalization theorem

Euler's generalization theorem states that, for a and n (with a < n) such that gcd(a, n) = 1:

$$a^{\phi(n)} \mod n = 1$$

To show that $a^{\phi(n)} \mod n = 1$, consider the reduced set of residues $\{r_i\}_{i=1,\dots,\phi(n)}$ and the (permuted) set of residues $\{r_j\}$:

$$\{r_j\} = \{ar_i \mod n\}_{i=1,...,\phi(n)} \{r_j\} = P \circ \{r_i\}_{i=1,...,\phi(n)}$$

Then the product of all the elements from the two reduced sets of residues, namely $\{r_i\}$ and $\{r_j\}$, must be equal:

$$\prod_{i=1}^{\phi(n)} r_i = \prod_{j=1}^{\phi(n)} r_j$$

Since the right-hand and left-hand sides of the equation are equal they should also be congruent modulo n:

$$\prod_{j=1}^{\phi(n)} r_j \equiv \prod_{i=1}^{\phi(n)} r_i \pmod{n}$$
$$\prod_{i=1}^{\phi(n)} (ar_i \mod n) \equiv \prod_{i=1}^{\phi(n)} r_i \pmod{n}$$
$$\prod_{i=1}^{\phi(n)} ar_i \equiv \prod_{i=1}^{\phi(n)} r_i \pmod{n}$$
$$a^{\phi(n)} \prod_{i=1}^{\phi(n)} r_i \equiv \prod_{i=1}^{\phi(n)} r_i \pmod{n}$$

because of the reducibility property. Dividing both sides by the factor $\prod_{i=1}^{\phi(n)} r_i$ leads to:

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

and since $1 \in \{0, \ldots, n-1\}$ then:

$$a^{\phi(n)} \mod n = 1$$

6.2 Fermat's little theorem

Fermat's little theorem states that if n is a prime number, with a < n, then:

$$a^{n-1} \mod n = 1$$

by property of the Euler function of a prime number, i.e. $\phi(n) = n - 1$.

6.3 Multiplicative inverses

Consider the expression

$$ax \mod n = 1$$

What is the multiplicative inverse x of a modulo n (assuming that gcd(a, n) = 1)? By Euler's generalization theorem:

 $ax \mod n = a^{\phi(n)} \mod n = 1$

which implies that:

$$x = a^{\phi(n)-1} \bmod n$$

Hence to compute an inverse a modular exponentiation program with the arguments $(a, [\phi(n) - 1], n)$ can be used. If n is a prime number, then $\phi(n) = n - 1$ (Fermat's theorem) and:

$$x = a^{(n-1)-1} \mod n = a^{n-2} \mod n$$

7 Galois Fields of Order p

Definition (Galois Field of Order p):

Let p be a prime number and $\mathbf{Z}_p = \{0, 1, \dots, p-1\}$ be the set of residues modulo p. The finite (Galois) field GF(p) is defined as the set \mathbf{Z}_p with the arithmetics modulo p.

Example(Galois Field modulo p = 5):

Consider the Galois Field of order p = 5, i.e. GF(5). Since p = 5 is a prime, the Galois field GF(5) consists of $\mathbb{Z}_5 = \{0, 1, 2, 3, 4\}$. The addition and multiplication operations in GF(5) are given in Table 1 as well as the additive and multiplicative inverses, -w and w^{-1} .

Table 1: Addition and multiplication operations in GF(5).

Addition					Multiplication							Inverses			
+	0	1	2	3	4	×	0	1	2	3	4		w	-w	w^{-1}
0	0	1	2	3	4	0	0	0	0	0	0		0	0	
1	1	2	3	4	0	1	0	1	2	3	4		1	4	1
2	2	3	4	0	1	2	0	2	4	1	3		2	3	3
3	3	4	0	1	2	3	0	3	1	4	2		3	2	2
4	4	0	1	2	3	4	0	4	3	2	1		4	1	4