

Covering Arrays and Extremal Set-Partition Systems

Lucia Moura School of Information Technology and Engineering University of Ottawa lucia@site.uottawa.ca

Workshop on Covering Arrays, May 2006

Definition: Covering Array

A covering array of strength t, k factors, g levels for each factor and size N, denoted by CA(N; t, k, g), is an $k \times N$ array with symbols from a g-ary alphabet G such that in every $t \times N$ subarray, every t-tuple in G^t is covered at least once.

$$N = 10, t = 2, k = 4, g = 3$$

1	2	3	4	5	6	7	8	9	10
0	0	0	1	1	1	2	2	2	0
0	1	2	0	1	2	0	1	2	0
0	1	2	2	0	1	1	2	0	1
0	0	2	1	2	0	2	0	1	1

Covering array optimization questions

Fix t and g.

Minimizing N for fixed k (number of tests)

 $CAN(t, k, g) = \min\{N : there \ exists \ a \ CA(N; t, k, g)\}.$

Maximizing k for fixed N (number of factors)

 $CAK(N, t, g) = \max\{k : there \ exists \ a \ CA(N; t, k, g)\}.$

Relationship between min-max problems

$$CAN(t,k,g) = \min\{N : CAK(N,t,g) \ge k\}.$$

We will focus on the following problem:

Maximizing k for fixed N (number of factors)

 $CAK(N, t, g) = \max\{k : there \ exists \ a \ CA(N; t, k, g)\}.$

General methodology for an optimization problem (maximization)

- Relaxed problem: relax constraints to find upper bounds
- Hard solution matching upper bound: build a feasible solution to hard problem that matches the upper bound (or close to upper bound)

More structure: sometimes it is worth adding more structure to the objects saught.

Final remarks

Binary covering arrays and set systems

 A binary covering array (strength 2) is a set system

Covering Arrays and Extremal Set-Partition Systems

۲

Binary covering arrays and set systems

 A binary covering array (strength 2) is a set system



{2,4,6} {2,4,5} {1,4,6}

Covering Arrays and Extremal Set-Partition Systems

Binary covering arrays and set systems

 A binary covering array (strength 2) is a set system



• Maximization problem: $CAK(N,t,g) = \max\{k : there \ exists \ a \ CA(N;t,k,g)\}$ ۲

Final remarks

Binary covering arrays and set systems

 A binary covering array (strength 2) is a set system

1	2	3	4	5	6	
0	0	1	1	1	0	{3,4,5 }
0	1	0	1	0	1	{2,4,6}
0	1	0	1	1	0	{2,4,5}
1	0	0	1	0	1	{1,4,6}

• Maximization problem: $CAK(N,t,g) = \max\{k : there \ exists \ a \ CA(N;t,k,g)\}$

• Maximization problem: strongly intersecting Given N, find a set system \mathcal{A} with maximum $|\mathcal{A}|$ such that for all $A, B \in \mathcal{A}$ we have:

$$\begin{split} &A\cap B\neq \emptyset, \qquad A\cap \overline{B}\neq \emptyset, \\ &\overline{A}\cap B\neq \emptyset, \qquad \overline{A}\cap \overline{B}\neq \emptyset. \end{split}$$

Sperner theorem for set systems

A system of subsets of an *n*-set has the *Sperner property* if no two subsets in the system are comparable.



Sperner's Theorem (1928)

If \mathcal{A} has the Sperner property, then $|\mathcal{A}| \leq {n \choose \lfloor \frac{n}{2} \rfloor}$.

The upper bound is only acchieved by the set of all $\left(\lfloor \frac{n}{2} \rfloor\right)$ -subsets of the *n*-set, or by its (subsetwise) complement.

A system of subsets of an *n*-set is *t*-intersecting if every two subsets in the system have intersection cardinality at least *t*. Example: n = 6, k = 3(t = 2) $\mathcal{A} = \{\{1, 2, 3\}, \{1, 2, 4\}, \{1, 3, 4\}, \{2, 3, 4\}\}$ (t = 1) $\mathcal{B} = \{\{1, 2, 3\}, \{1, 2, 4\}, \{1, 2, 5\}, \{1, 2, 6\}, \{1, 3, 4\}, \{1, 3, 5\}, \{1, 3, 6\}, \{1, 4, 5\}, \{1, 4, 6\}, \{1, 5, 6\}\}$

Erdos-Ko-Rado Theorem (1961)

Let t < k < n. Let \mathcal{A} be a *t*-intersecting system of subsets of an *n*-set, such that each subset has cardinality at most *k*. If $n \ge (t+1)(k-t+1)$, then $|\mathcal{A}| \le {n-t \choose k-t}$.

Moreover, if n > (t + 1)(k - t + 1), then equality holds if and only if A is a k-uniform trivially t-intersecting system.

- 4 同 ト 4 回 ト -

Applying the methodology to covering arrays

Methodology for binary covering arrays

- Relaxed problem: strongly intersecting set system (plus complement) is Sperner - use Sperner upper bound.
- e Hard solution matching upper bound: build a strongly intersecting set system that matches the upper bound.

More structure: use point-balanced covering arrays (uniform set systems with sets of cardinality around N/2).

Solving the binary covering array problem

Pick all $\lfloor n/2 \rfloor$ -subsets of $\lfloor 1, n \rfloor$ that contain a common element.

	n even.
	1 2 3 4 5 6
	111000
n odd:	110100
1 2 3 4 5	110010
11000	$1\ 1\ 0\ 0\ 1$
10100	101100
10010	101010
10001	$1\ 0\ 1\ 0\ 1$
	100110
	$1 \ 0 \ 0 \ 1 \ 0 \ 1$
	$1 \ 0 \ 0 \ 1 \ 1$

Note: Systems are strongly Sperner and 1-intersecting.

Theorem (Katona 1973, Kleitman and Spencer 1973)

 $CAK(n, t = 2, g = 2) = {\binom{n-1}{\lfloor n/2 \rfloor - 1}}$. Moreover, this bound is attained by a $\lfloor n/2 \rfloor$ -uniform trivially 1-intersecting set system.

Proof: Let \mathcal{A} be the set system corresponding to the CA.

• (Case 1) n even.

Theorem (Katona 1973, Kleitman and Spencer 1973)

 $CAK(n, t = 2, g = 2) = \binom{n-1}{\lfloor n/2 \rfloor - 1}$. Moreover, this bound is attained by a $\lfloor n/2 \rfloor$ -uniform trivially 1-intersecting set system.

- (Case 1) n even.
- $\mathcal{A}^* = \{A, \overline{A} : A \in \mathcal{A}\}$ is Sperner.

Theorem (Katona 1973, Kleitman and Spencer 1973)

 $CAK(n, t = 2, g = 2) = \binom{n-1}{\lfloor n/2 \rfloor - 1}$. Moreover, this bound is attained by a $\lfloor n/2 \rfloor$ -uniform trivially 1-intersecting set system.

- (Case 1) n even.
- $\mathcal{A}^* = \{A, \overline{A} : A \in \mathcal{A}\}$ is Sperner.
- Sperner's theorem implies $|\mathcal{A}^*| \leq {n \choose n/2}$.

Theorem (Katona 1973, Kleitman and Spencer 1973)

 $CAK(n, t = 2, g = 2) = \binom{n-1}{\lfloor n/2 \rfloor - 1}$. Moreover, this bound is attained by a $\lfloor n/2 \rfloor$ -uniform trivially 1-intersecting set system.

- (Case 1) n even.
- $\mathcal{A}^* = \{A, \overline{A} : A \in \mathcal{A}\}$ is Sperner.
- Sperner's theorem implies $|\mathcal{A}^*| \leq {n \choose n/2}$.
- $|\mathcal{A}| \leq \frac{1}{2}|\mathcal{A}^*| \leq \frac{1}{2}\binom{n}{n/2} = \binom{n-1}{n/2-1}.$

Theorem (Katona 1973, Kleitman and Spencer 1973)

 $CAK(n, t = 2, g = 2) = \binom{n-1}{\lfloor n/2 \rfloor - 1}$. Moreover, this bound is attained by a $\lfloor n/2 \rfloor$ -uniform trivially 1-intersecting set system.

- (Case 1) n even.
- $\mathcal{A}^* = \{A, \overline{A} : A \in \mathcal{A}\}$ is Sperner.
- Sperner's theorem implies $|\mathcal{A}^*| \leq {n \choose n/2}$.
- $|\mathcal{A}| \leq \frac{1}{2}|\mathcal{A}^*| \leq \frac{1}{2}\binom{n}{n/2} = \binom{n-1}{n/2-1}.$
- (Case 2) n odd.

Theorem (Katona 1973, Kleitman and Spencer 1973)

 $CAK(n, t = 2, g = 2) = \binom{n-1}{\lfloor n/2 \rfloor - 1}$. Moreover, this bound is attained by a $\lfloor n/2 \rfloor$ -uniform trivially 1-intersecting set system.

- (Case 1) n even.
- $\mathcal{A}^* = \{A, \overline{A} : A \in \mathcal{A}\}$ is Sperner.
- Sperner's theorem implies $|\mathcal{A}^*| \leq {n \choose n/2}$.
- $|\mathcal{A}| \leq \frac{1}{2}|\mathcal{A}^*| \leq \frac{1}{2}\binom{n}{n/2} = \binom{n-1}{n/2-1}.$
- (Case 2) n odd.
- Wlog assume $|A| \leq \lfloor n/2 \rfloor$, for all $A \in \mathcal{A}$.

Theorem (Katona 1973, Kleitman and Spencer 1973)

 $CAK(n, t = 2, g = 2) = \binom{n-1}{\lfloor n/2 \rfloor - 1}$. Moreover, this bound is attained by a $\lfloor n/2 \rfloor$ -uniform trivially 1-intersecting set system.

- (Case 1) n even.
- $\mathcal{A}^* = \{A, \overline{A} : A \in \mathcal{A}\}$ is Sperner.
- Sperner's theorem implies $|\mathcal{A}^*| \leq {n \choose n/2}$.
- $|\mathcal{A}| \leq \frac{1}{2}|\mathcal{A}^*| \leq \frac{1}{2}\binom{n}{n/2} = \binom{n-1}{n/2-1}.$
- (Case 2) n odd.
- Wlog assume $|A| \leq \lfloor n/2 \rfloor$, for all $A \in \mathcal{A}$.
- \mathcal{A} is 1-intersecting, so by the EKR theorem, $|\mathcal{A}| \leq {n-1 \choose \lfloor n/2 \rfloor 1}$.

Covering arrays are systems of set partitions

• A covering array (strength 2) is a system of set-partitions:

	2	3	4	5	0	1	0	9	10
0	0	0	1	1	1	2	2	2	0
0	1	2	0	1	2	0	1	2	0
0	1	2	2	0	1	1	2	0	1
0	0	2	1	2	0	2	0	1	1

{1,2,3,10} {4,5,6} {7,8,9} {1,4,7,10} {2,5,8} {3,6,9} {1,5,9} {2,6,7,10} {3,4,8} {1,2,6,8} {4,9,10} {3,5,7} Covering arrays are systems of set partitions

• A covering array (strength 2) is a system of set-partitions:

1	2	3	4	5	0	1	0	9	10
0	0	0	1	1	1	2	2	2	0
0	1	2	0	1	2	0	1	2	0
0	1	2	2	0	1	1	2	0	1
0	0	2	1	2	0	2	0	1	1

*{*1*,*2*,*3*,*10*} <i>{*4*,*5*,*6*} <i>{*7*,*8*,*9*} {*1,4,7,10*} {*2,5,8*} <i>{*3,6,9*} {1,5,9} <i>{2,6,7,10} <i>{3,4,8} {*1,2,6,8*} {***4,9,10***} {***3,5,7***}*

Maximization problem:

Given N, find a set partition system \mathcal{P} with maximum $|\mathcal{P}|$ that is (pairwise) strongly intersecting: For all $P, Q \in \mathcal{P}$ we have

$$\text{for all } P_i \in P, Q_j \in Q, \quad P_i \cap Q_j \neq \emptyset.$$

Strongly intersecting condition: upper bound via 2-parts

Theorem (Stevens, Moura and Mendelsohn 1998)

$$CAK(n,2,g) \leq \frac{1}{2} {\lfloor \frac{2n}{g} \rfloor \choose \lfloor \frac{n}{g} \rfloor}.$$

This theorem only uses the two smallest parts of each partition, and the following fact:

Consider a pair of set systems, A_1, A_2, \ldots, A_k and B_1, B_2, \ldots, B_k , with $|A_i| + |B_i| \le c$ and $|A_i| \le a \le c/2$, and such that $A_i \cap B_i = \emptyset$, and all other sets intersect. Then, $k \le \frac{1}{2} \binom{c}{a}$. It is possible to relabel symbols of the covering array so that $|P_{1j}| \le \lfloor \frac{n}{g} \rfloor$ and $|P_{1j}| + |P_{2j}| \le \lfloor \frac{2n}{g} \rfloor$

Stronly intersecting versus Sperner formulation

Strongly intersecting formulation:

Partitions P and Q corresponding to two rows of a covering array must satisfy:

for all
$$P_i \in P, Q_j \in Q, \quad P_i \cap Q_j \neq \emptyset.$$

Strongly Sperner formulation:

Partitions P and Q corresponding to two rows of a covering array must satisfy:

for all
$$P_i \in P, Q_j \in Q, \quad P_i \not\subseteq \overline{Q}_j$$
 and $Q_j \not\subseteq \overline{P}_i$

 Framework by Ahlswede, Cai and Zhang: System of "clouds": A = {A₁, A₂,..., A_k}.

Covering Arrays and Extremal Set-Partition Systems

- Framework by Ahlswede, Cai and Zhang: System of "clouds": $\mathcal{A} = \{\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k\}$.
 - Types: (\forall, \forall) , (\exists, \exists) , (\forall, \exists) , (\exists, \forall)

- Framework by Ahlswede, Cai and Zhang: System of "clouds": $\mathcal{A} = \{\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k\}.$
 - Types: (\forall, \forall) , (\exists, \exists) , (\forall, \exists) , (\exists, \forall)
 - Binary relations: Comparable, iNcomparable, Disjoint, Intersecting.

- Framework by Ahlswede, Cai and Zhang: System of "clouds": $\mathcal{A} = \{\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k\}.$
 - Types: (\forall, \forall) , (\exists, \exists) , (\forall, \exists) , (\exists, \forall)
 - Binary relations: Comparable, iNcomparable, Disjoint, Intersecting.
 - $I_n(\exists, \forall)$: largest cardinality k of a system of clouds of [1, n] such that for all $A_i, A_j, \in A$:

 $\exists A \in \mathcal{A}_i, \forall A' \in \mathcal{A}_j (A \cap A' \neq \emptyset).$

- Framework by Ahlswede, Cai and Zhang: System of "clouds": A = {A₁, A₂,..., A_k}.
 - Types: (\forall, \forall) , (\exists, \exists) , (\forall, \exists) , (\exists, \forall)
 - Binary relations: Comparable, iNcomparable, Disjoint, Intersecting.
 - $I_n(\exists, \forall)$: largest cardinality k of a system of clouds of [1, n] such that for all $A_i, A_j, \in A$:

 $\exists A \in \mathcal{A}_i, \forall A' \in \mathcal{A}_j (A \cap A' \neq \emptyset).$

• For us each cloud is a g-partition of [1, n]:

- Framework by Ahlswede, Cai and Zhang: System of "clouds": A = {A₁, A₂,..., A_k}.
 - Types: (\forall, \forall) , (\exists, \exists) , (\forall, \exists) , (\exists, \forall)
 - Binary relations: Comparable, iNcomparable, Disjoint, Intersecting.
 - $I_n(\exists, \forall)$: largest cardinality k of a system of clouds of [1, n] such that for all $A_i, A_j, \in A$:

 $\exists A \in \mathcal{A}_i, \forall A' \in \mathcal{A}_j (A \cap A' \neq \emptyset).$

- For us each cloud is a g-partition of [1, n]:
 - $I_n(\forall,\forall) = CAK(n,2,g).$

- Framework by Ahlswede, Cai and Zhang: System of "clouds": A = {A₁, A₂,..., A_k}.
 - Types: (\forall, \forall) , (\exists, \exists) , (\forall, \exists) , (\exists, \forall)
 - Binary relations: Comparable, iNcomparable, Disjoint, Intersecting.
 - $I_n(\exists, \forall)$: largest cardinality k of a system of clouds of [1, n] such that for all $A_i, A_j, \in A$:

 $\exists A \in \mathcal{A}_i, \forall A' \in \mathcal{A}_j (A \cap A' \neq \emptyset).$

- For us each cloud is a g-partition of [1, n]:
 - $I_n(\forall,\forall) = CAK(n,2,g).$
 - $N_n(\forall, \forall) = \text{Sperner } g\text{-partition system}$ Note: $I_n(\forall, \forall) \le N_n(\forall, \forall)$

Sperner's theorem for set-partition systems

 $N_n(\forall, \forall)$: largest cardinality k of a system \mathcal{P} of g-partitions of [1, n] such that for all $\mathcal{P}_i, \mathcal{P}_j \in \mathcal{P}$:

 $\forall P \in \mathcal{P}_i, \forall P' \in \mathcal{P}_j(P \not\subseteq P' \text{and} P' \not\subseteq P).$ (Weakly) Sperner

Theorem (Meagher, Moura and Stevens 2005)

Let g, n such that n = cg + r and $0 \le r < g$. Then,

$$N_n(orall,orall) \leq rac{1}{(g-r)+rac{r(c+1)}{n-1}} inom{n}{c}.$$

Theorem (Meagher, Moura and Stevens 2005)

Let g, n such that g|n. Then, $N_n(\forall, \forall) = \binom{n-1}{\frac{n}{g}-1}$. Moreover, this bound is met if and only if the g-partitions are uniform (all parts with cardinality $\frac{n}{g}$).

Example: weakly Sperner property

n=2g

{1,2,3},{4,5,6}

{1,2,4},{3,5,6}

{1,2,5},{3,4,6}

{1,2,6},{3,4,5}

{1,3,4},{2,5,6}

{1,3,5},{2,4,6}

{1,3,6},{2,4,5}

{1,4,5},{2,3,6} {1,4,6},{2,3,5}

{1,5,6},{2,3,4}

n=3g

{1,2,3},{4,5,6},{7,8,9} {1,2,4},.... {1,2,5},... {1.2.6}....

{1,7,8},... {1,8,9},{2,3,4}, <mark>{5,6,7}</mark>

-

Comparison of two bounds obtained

Theorem (Stevens, Moura and Mendelsohn 1998)

 $CAK(n,2,g) \leq \frac{1}{2} {\binom{\lfloor \frac{2n}{g} \rfloor}{\lfloor \frac{n}{g} \rfloor}}.$

Theorem (Meagher, Moura and Stevens 2005)

If g|n, then $CAK(n, 2, g) \leq {\binom{n-1}{\frac{n}{g}-1}}$.

if g > 2, g|n, then

$$\frac{1}{2} \binom{\frac{2n}{g}}{\frac{n}{g}} < \binom{n-1}{\frac{n}{g}-1}$$

Covering Arrays and Extremal Set-Partition Systems

Erdos-Ko-Rado theorem for set-partition systems

- We are interested on: (\exists, \exists) with property *p*-intersecting.
- This is useful for bounds on "anti-covering-arrays" for certain uniform cases. For example: $n = g^2$, p = 2

Conjecture

Suppose g|n, and let c = n/g be the size of each part of the (uniform) partition system. $p - I_n(\exists, \exists) = \binom{n-p}{c-p}U(n-c, g-1)$.

This has been proven for p = c:

Theorem (Meagher and Moura 2005)

Let $n \ge g \ge 1$ and let $\mathcal{P} \subseteq U_g^n$ be a partition system in which every two partitions share at least one class. Let c = n/g. Then, $|\mathcal{P}| \le U(n - c, g - 1)$

∃ ► < ∃ ►</p>

Final remarks

Higher strength: strongly intersecting and Sperner

Strongly intersecting formulation:

Partitions $P^{i_1}, P^{i_2}, \ldots, P^{i_t}$ corresponding to t rows of a covering array must satisfy:

for all
$$A_{k_1} \in P^{i_1}, A_{k_2} \in P^{i_2}, \dots, A_{k_t} \in P^{i_t},$$

$$A_{k_1} \cap A_{k_2} \cap \dots \cap A_{k_t} \neq \emptyset.$$

Generalization of *t*-wise intersecting set systems.

Strongly Sperner formulation:

Partitions $P^{i_1}, P^{i_2}, \dots P^{i_t}$ corresponding to t rows of a covering array must satisfy:

for all
$$A_{k_1} \in P^{i_1}, A_{k_2} \in P^{i_2}, \dots, A_{k_t} \in P^{i_t},$$

 $A_{k_1} \not\subseteq \overline{A}_{k_2} \cup \dots \cup \overline{A}_{k_t},$

- Binary case was solved via systems of sets.
 The case g > 2 can be studied via systems of g-partition:
 - Stronger upper bounds for set-partition systems= stronger lower bounds for covering arrays.
 - Some cases might be amenable to complete solution (for example, binary covering arrays with strength 3.)
- Special attention should be given to point-balanced case (Meagher conjectures there is always an optimal covering array that is (almost) point-balanced).
- Systems of set-partitions are interesting on their own right, and other extremal problems could be investigated.

< 同 ▶

∃ → < ∃ →</p>

References

- R. Ahlswede, N. Cai and Z. Zhang. Higher level extremal problems. J. Combin. Inform. System Sci., 21:185–210, 1996.
- P. Erdos, C. Ko and R. Rado. Intersection theorems for systems of finite sets. *Quart. J. Math. Oxford Ser.* 12:313–320, 1961.
- G. Katona. Two applications (for search theory and truth functions) of Sperner type theorems. Period. Math. Hungar., 3:19–26, 1973.
- 4 D.J. Kleitman and J. Spencer. Families of k-independent sets. Discrete Math., 6:255–262, 1973.
- K. Meagher and L. Moura. Erdos-Ko-Rado theorems for uniform set-partition systems. Electr. J. Combin. 12, R40, 12 pages, 2005.
- K. Meagher, L. Moura and B. Stevens. A Sperner-type theorem for set-partition systems. *Electr. J. Combin.* 12, N20, 6 pages, 2005.
- 🚺 E. Sperner. Ein Satz uber Untermengen einer endlichen Menge. Math. Z., 27:544-548, 1928.
- B. Stevens, L. Moura and E. Mendelsohn. Lower bounds for transversal covers. Des., Codes and Cryptogr., 15:279-299, 1998.