# Covering Arrays and Extremal Set-Partition Systems 

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## Covering arrays

## Definition: Covering Array

A covering array of strength $t, k$ factors, $g$ levels for each factor and size $N$, denoted by $C A(N ; t, k, g)$, is an $k \times N$ array with symbols from a $g$-ary alphabet $G$ such that in every $t \times N$ subarray, every $t$-tuple in $G^{t}$ is covered at least once.

$$
N=10, t=2, k=4, g=3
$$

| 1 | 2 | 3 | $\mathbf{4}$ | 5 | 6 | $\mathbf{7}$ | 8 | $\mathbf{9}$ | 10 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 1 | 1 | 1 | 2 | 2 | 2 | 0 |
| 0 | 1 | 2 | 0 | 1 | 2 | 0 | 1 | 2 | 0 |
| 0 | 1 | 2 | 2 | 0 | 1 | 1 | 2 | 0 | 1 |
| 0 | 0 | 2 | 1 | 2 | 0 | 2 | 0 | 1 | 1 |

## Covering array optimization questions

Fix $t$ and $g$.
Minimizing $N$ for fixed $k$ (number of tests)

$$
C A N(t, k, g)=\min \{N: \text { there exists a } C A(N ; t, k, g)\} .
$$

Maximizing $k$ for fixed $N$ (number of factors)

$$
C A K(N, t, g)=\max \{k: \text { there exists a } C A(N ; t, k, g)\} .
$$

Relationship between min-max problems

$$
C A N(t, k, g)=\min \{N: C A K(N, t, g) \geq k\}
$$

## Methodology

We will focus on the following problem:
Maximizing $k$ for fixed $N$ (number of factors)

$$
C A K(N, t, g)=\max \{k: \text { there exists a } C A(N ; t, k, g)\} .
$$

General methodology for an optimization problem (maximization)
(1) Relaxed problem: relax constraints to find upper bounds
(2) Hard solution matching upper bound: build a feasible solution to hard problem that matches the upper bound (or close to upper bound)

More structure: sometimes it is worth adding more structure to the objects saught.

## Binary covering arrays and set systems

- A binary covering array (strength 2) is a set system


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| 1 | 2 | 3 | 4 | 5 | 6 |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 1 | 1 | 0 | $\{3,4,5\}$ |
| 0 | 1 | 0 | 1 | 0 | 1 | $\{2,4,6\}$ <br> 0 1 |
|  | 0 | 1 | 1 | 0 | $\{2,4,5\}$ |  |
| 1 | 0 | 0 | 1 | 0 | 1 | $\{1,4,6\}$ |

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\{2,4,6\}
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| 1 |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 1 | 1 | 0 |
| 0 | 1 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 | 0 | 1 |

- Maximization problem:
$\operatorname{CAK}(N, t, g)=\max \{k$ : there exists a $C A(N ; t, k, g)\}$
- Maximization problem: strongly intersecting
Given $N$, find a set system $\mathcal{A}$ with maximum $|\mathcal{A}|$ such that for all $A, B \in \mathcal{A}$ we have:

$$
\begin{array}{ll}
A \cap B \neq \emptyset, & A \cap \bar{B} \neq \emptyset \\
\bar{A} \cap B \neq \emptyset, & \bar{A} \cap \bar{B} \neq \emptyset .
\end{array}
$$

## Sperner theorem for set systems

A system of subsets of an $n$-set has the Sperner property if no two subsets in the system are comparable.


## Sperner's Theorem (1928)

If $\mathcal{A}$ has the Sperner property, then $|\mathcal{A}| \leq\binom{ n}{\left\lfloor\frac{n}{2}\right\rfloor}$.
The upper bound is only acchieved by the set of all $\left(\left\lfloor\frac{n}{2}\right\rfloor\right)$-subsets of the $n$-set, or by its (subsetwise) complement.

## Erdos-Ko-Rado theorem for set systems

A system of subsets of an $n$-set is $t$-intersecting if every two subsets in the system have intersection cardinality at least $t$.
Example: $n=6, k=3$
$(t=2) \quad \mathcal{A}=\{\{1,2,3\},\{1,2,4\},\{1,3,4\},\{2,3,4\}\}$
$(t=1) \quad \mathcal{B}=\{\{1,2,3\},\{1,2,4\},\{1,2,5\},\{1,2,6\},\{1,3,4\}$,
$\{1,3,5\},\{1,3,6\},\{1,4,5\},\{1,4,6\},\{1,5,6\}\}$

## Erdos-Ko-Rado Theorem (1961)

Let $t<k<n$. Let $\mathcal{A}$ be a $t$-intersecting system of subsets of an $n$-set, such that each subset has cardinality at most $k$.
If $n \geq(t+1)(k-t+1)$, then $|\mathcal{A}| \leq\binom{ n-t}{k-t}$.
Moreover, if $n>(t+1)(k-t+1)$, then equality holds if and only if $\mathcal{A}$ is a $k$-uniform trivially $t$-intersecting system.

## Applying the methodology to covering arrays

## Methodology for binary covering arrays

(1) Relaxed problem: strongly intersecting set system (plus complement) is Sperner - use Sperner upper bound.
(2) Hard solution matching upper bound: build a strongly intersecting set system that matches the upper bound.

More structure: use point-balanced covering arrays (uniform set systems with sets of cardinality around $N / 2$ ).

## Solving the binary covering array problem

Pick all $\lfloor n / 2\rfloor$-subsets of $[1, n]$ that contain a common element.


Note: Systems are strongly Sperner and 1-intersecting.

## The binary covering array theorem

Theorem (Katona 1973, Kleitman and Spencer 1973)
$C A K(n, t=2, g=2)=\binom{n-1}{\lfloor n / 2\rfloor-1}$. Moreover, this bound is attained by a $\lfloor n / 2\rfloor$-uniform trivially 1 -intersecting set system.

Proof: Let $\mathcal{A}$ be the set system corresponding to the CA.

- (Case 1) $n$ even.


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- $|\mathcal{A}| \leq \frac{1}{2}\left|\mathcal{A}^{*}\right| \leq \frac{1}{2}\binom{n}{n / 2}=\binom{n-1}{n / 2-1}$. $\square$


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- Wlog assume $|A| \leq\lfloor n / 2\rfloor$, for all $A \in \mathcal{A}$.


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- (Case 1) $n$ even.
- $\mathcal{A}^{*}=\{A, \bar{A}: A \in \mathcal{A}\}$ is Sperner.
- Sperner's theorem implies $\left|\mathcal{A}^{*}\right| \leq\binom{ n}{n / 2}$.
- $|\mathcal{A}| \leq \frac{1}{2}\left|\mathcal{A}^{*}\right| \leq \frac{1}{2}\binom{n}{n / 2}=\binom{n-1}{n / 2-1}$. $\square$
- (Case 2) $n$ odd.
- Wlog assume $|A| \leq\lfloor n / 2\rfloor$, for all $A \in \mathcal{A}$.
- $\mathcal{A}$ is 1 -intersecting, so by the EKR theorem, $|\mathcal{A}| \leq\binom{ n-1}{\lfloor n / 2\rfloor-1}$.


## Covering arrays are systems of set partitions

- A covering array (strength 2 ) is a system of set-partitions:

| $\begin{array}{llllllllll}3 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10\end{array}$ |  |  |  |  |  |  |  |  |  | \{1,2,3,10\} \{4,5,6\} \{7,8,9\} |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 1 | 1 | 2 | 2 | 2 | 0 |  |  |
| 0 | 1 | 2 | 0 | 1 | 2 | 0 | 1 | 2 | 0 | \{1,4,7,10\} \{2,5,8\} | \{3,6,9\} |
| 0 | 1 | 2 | 2 | 0 | 1 | 1 | 2 | 0 | 1 | $\{1,5,9\}\{2,6,7,10\}$ | \{3,4,8\} |
| 0 | 0 | 2 | 1 | 2 | 0 | 2 | 0 | 1 | 1 | \{1,2,6,8\} \{4,9,10\} | \{3,5,7\} |

## Covering arrays are systems of set partitions

- A covering array (strength 2 ) is a system of set-partitions:

- Maximization problem:

Given $N$, find a set partition system $\mathcal{P}$ with maximum $|\mathcal{P}|$ that is (pairwise) strongly intersecting:
For all $P, Q \in \mathcal{P}$ we have

$$
\text { for all } P_{i} \in P, Q_{j} \in Q, \quad P_{i} \cap Q_{j} \neq \emptyset
$$

## Strongly intersecting condition: upper bound via 2-parts

Theorem (Stevens, Moura and Mendelsohn 1998)
$C A K(n, 2, g) \leq \frac{1}{2}\left(\begin{array}{c}\left\lfloor\begin{array}{l}\left.\frac{2 n}{g}\right\rfloor \\ \left\lfloor\frac{n}{g}\right. \\ \hline\end{array}\right)\end{array}\right)$.
This theorem only uses the two smallest parts of each partition, and the following fact:
Consider a pair of set systems, $A_{1}, A_{2}, \ldots, A_{k}$ and $B_{1}, B_{2}, \ldots, B_{k}$, with $\left|A_{i}\right|+\left|B_{i}\right| \leq c$ and $\left|A_{i}\right| \leq a \leq c / 2$, and such that $A_{i} \cap B_{i}=\emptyset$, and all other sets intersect. Then, $k \leq \frac{1}{2}\binom{c}{a}$.
It is possible to relabel symbols of the covering array so that $\left|P_{1 j}\right| \leq\left\lfloor\frac{n}{g}\right\rfloor$ and $\left|P_{1 j}\right|+\left|P_{2 j}\right| \leq\left\lfloor\frac{2 n}{g}\right\rfloor$

## Stronly intersecting versus Sperner formulation

## Strongly intersecting formulation:

Partitions $P$ and $Q$ corresponding to two rows of a covering array must satisfy:

$$
\text { for all } P_{i} \in P, Q_{j} \in Q, \quad P_{i} \cap Q_{j} \neq \emptyset .
$$

## Strongly Sperner formulation:

Partitions $P$ and $Q$ corresponding to two rows of a covering array must satisfy:

$$
\text { for all } P_{i} \in P, Q_{j} \in Q, \quad P_{i} \nsubseteq \bar{Q}_{j} \text { and } Q_{j} \nsubseteq \bar{P}_{i}
$$

## Higher level extremal problems

- Framework by Ahlswede, Cai and Zhang: System of "clouds": $\mathcal{A}=\left\{\mathcal{A}_{1}, \mathcal{A}_{2}, \ldots, \mathcal{A}_{k}\right\}$.


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- Binary relations: Comparable, iNcomparable, Disjoint, Intersecting.


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Comparable, iNcomparable, Disjoint, Intersecting.

- $I_{n}(\exists, \forall)$ : largest cardinality $k$ of a system of clouds of $[1, n]$ such that for all $\mathcal{A}_{i}, \mathcal{A}_{j}, \in \mathcal{A}$ :

$$
\exists A \in \mathcal{A}_{i}, \forall A^{\prime} \in \mathcal{A}_{j}\left(A \cap A^{\prime} \neq \emptyset\right)
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- For us each cloud is a $g$-partition of $[1, n]$ :


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$$
\text { - } I_{n}(\forall, \forall)=C A K(n, 2, g) \text {. }
$$

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$$
\exists A \in \mathcal{A}_{i}, \forall A^{\prime} \in \mathcal{A}_{j}\left(A \cap A^{\prime} \neq \emptyset\right)
$$

- For us each cloud is a $g$-partition of $[1, n]$ :
- $I_{n}(\forall, \forall)=C A K(n, 2, g)$.
- $N_{n}(\forall, \forall)=$ Sperner $g$-partition system

Note: $I_{n}(\forall, \forall) \leq N_{n}(\forall, \forall)$

## Sperner's theorem for set-partition systems

$N_{n}(\forall, \forall)$ : largest cardinality $k$ of a system $\mathcal{P}$ of $g$-partitions of $[1, n]$ such that for all $\mathcal{P}_{i}, \mathcal{P}_{j} \in \mathcal{P}$ :

$$
\forall P \in \mathcal{P}_{i}, \forall P^{\prime} \in \mathcal{P}_{j}\left(P \nsubseteq P^{\prime} \text { and } P^{\prime} \nsubseteq P\right) \text {. (Weakly) Sperner }
$$

Theorem (Meagher, Moura and Stevens 2005)
Let $g, n$ such that $n=c g+r$ and $0 \leq r<g$. Then,

$$
N_{n}(\forall, \forall) \leq \frac{1}{(g-r)+\frac{r(c+1)}{n-1}}\binom{n}{c}
$$

Theorem (Meagher, Moura and Stevens 2005)
Let $g, n$ such that $g \mid n$. Then, $N_{n}(\forall, \forall)=\binom{n-1}{\frac{n}{g}-1}$. Moreover, this bound is met if and only if the $g$-partitions are uniform (all parts with cardinality $\frac{n}{g}$ ).

## Example: weakly Sperner property

$$
n=2 g
$$

\{1,2,3\},\{4,5,6\}
\{1,2,4\},\{3,5,6\}
\{1,2,5\},\{3,4,6\}
\{1,2,6\},\{3,4,5\}
\{1,3,4\},\{2,5,6\}
\{1,3,5\},\{2,4,6\}
\{1,3,6\},\{2,4,5\}
\{1,4,5\},\{2,3,6\}
\{1,4,6\},\{2,3,5\}
\{1,5,6\},\{2,3,4\}
$n=3 g$
\{1,2,3\},\{4,5,6\},\{7,8,9\}
\{1,2,4\},....
\{1,2,5\},...
\{1,2,6\},...
\{1,7,8\},...
\{1,8,9\},\{2,3,4\}, \{5,6,7\}

## Comparison of two bounds obtained

Theorem (Stevens, Moura and Mendelsohn 1998)
$C A K(n, 2, g) \leq \frac{1}{2}\left(\begin{array}{c}\left\lfloor\begin{array}{c}\left.\frac{2 n}{g}\right\rfloor \\ \left\lfloor\frac{n}{g}\right. \\ \hline\end{array}\right)\end{array}\right)$.

Theorem (Meagher, Moura and Stevens 2005)
If $g \mid n$, then $\operatorname{CAK}(n, 2, g) \leq\binom{ n-1}{\frac{n}{g}-1}$.
if $g>2, g \mid n$, then

$$
\frac{1}{2}\binom{\frac{2 n}{g}}{\frac{n}{g}}<\binom{n-1}{\frac{n}{g}-1}
$$

## Erdos-Ko-Rado theorem for set-partition systems

- We are interested on: $(\exists, \exists)$ with property $p$-intersecting.
- This is useful for bounds on "anti-covering-arrays" for certain uniform cases. For example: $n=g^{2}, p=2$


## Conjecture

Suppose $g \mid n$, and let $c=n / g$ be the size of each part of the (uniform) partition system. $p-I_{n}(\exists, \exists)=\binom{n-p}{c-p} U(n-c, g-1)$.

This has been proven for $p=c$ :
Theorem (Meagher and Moura 2005)
Let $n \geq g \geq 1$ and let $\mathcal{P} \subseteq U_{g}^{n}$ be a partition system in which every two partitions share at least one class. Let $c=n / g$. Then, $|\mathcal{P}| \leq U(n-c, g-1)$

## Higher strength: strongly intersecting and Sperner

## Strongly intersecting formulation:

Partitions $P^{i_{1}}, P^{i_{2}}, \ldots, P^{i_{t}}$ corresponding to $t$ rows of a covering array must satisfy:

$$
\text { for all } \begin{array}{r}
A_{k_{1}} \in P^{i_{1}}, A_{k_{2}} \in P^{i_{2}}, \ldots, A_{k_{t}} \in P^{i_{t}}, \\
\\
A_{k_{1}} \cap A_{k_{2}} \cap \cdots \cap A_{k_{t}} \neq \emptyset .
\end{array}
$$

Generalization of $t$-wise intersecting set systems.

## Strongly Sperner formulation:

Partitions $P^{i_{1}}, P^{i_{2}}, \ldots P^{i_{t}}$ corresponding to $t$ rows of a covering array must satisfy:

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\text { for all } \begin{array}{r}
A_{k_{1}} \in P^{i_{1}}, A_{k_{2}} \in P^{i_{2}}, \ldots, A_{k_{t}} \in P^{i_{t}}, \\
A_{k_{1}} \nsubseteq \bar{A}_{k_{2}} \cup \cdots \cup \bar{A}_{k_{t}},
\end{array}
$$

## Conclusions

- Binary case was solved via systems of sets.

The case $g>2$ can be studied via systems of $g$-partition:

- Stronger upper bounds for set-partition systems= stronger lower bounds for covering arrays.
- Some cases might be amenable to complete solution ( for example, binary covering arrays with strength 3.)
- Special attention should be given to point-balanced case (Meagher conjectures there is always an optimal covering array that is (almost) point-balanced).
- Systems of set-partitions are interesting on their own right, and other extremal problems could be investigated.


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