## ELG-5373 Secure Communications and Data Encryption

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Assignment #2 (due on Wednesday, February 27, 2002 at the beginning of the lecture.)
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## Question 1:

(permutation)
Let $\Pi(m)$ be a permutation of the $n$-bit integers: $0,1, \ldots, 2^{n}-1$, where $0 \leq m \leq 2^{n}$. For instance, the standard permutation $P$ in the DES algorithm is a permutation for 32-bit integers. If $\prod(m)=m$ then this value of $m$ is called a fixed point in the permutation.
a) Find an expression for the probability $P_{\text {no fixed point }}$ as a function of $N=2^{n}$. Hint: Consider the set of permutations $S_{N}$ on $\left[0, \ldots, 2^{n}-1\right]$.
b) Show that more than $60 \%$ of the mappings will have at least one fixed point.

## Question 2:

(DES modes of operation)
Problem 3.4 from the course notes.

## Question 3:

(double DES)
Consider double DES encryption with keys $K_{1}$ and $K_{2}: C=D E S_{K_{2}}\left[D E S_{K_{1}}(M)\right]$. If $D E S_{K_{2}}(X)=$ $D E S_{K_{1}}^{-1}(X)$, then $K_{1}$ and $K_{2}$ are called dual keys. This undesirable since the ciphertext $C$ will be the original plaintext $M$. Now a key $K$ will be a self-dual key if it is its own dual key.
a) Show that if $C_{0}$ is either all 0 's or all 1 's and $D_{0}$ is either all 0 's or all 1 's, then the key $K$ is a self-dual key.
b) Show that the following keys (in hexadecimal form) are self-dual:

$$
\begin{array}{lllllllllllllllll}
K_{1}= & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \\
K_{2}= & F & E & F & E & F & E & F & E & F & E & F & E & F & E & F & E \\
K_{3}= & 1 & F & 1 & F & 1 & F & 1 & F & 0 & E & 0 & E & 0 & E & 0 & E \\
K_{4}= & E & 0 & E & 0 & E & 0 & E & 0 & F & 1 & F & 1 & F & 1 & F & 1
\end{array}
$$

c) Show that the following pairs of keys are dual:

| $K_{1,1}=$ | $E$ | 0 | 0 | 1 | $E$ | 0 | 0 | 1 | $F$ | 1 | 0 | 1 | $F$ | 1 | 0 | 1 |
| ---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| and $K_{1,2}=$ | 0 | 1 | $E$ | 0 | 0 | 1 | $E$ | 0 | 0 | 1 | $F$ | 1 | 0 | 1 | $F$ | 1 |
| $K_{2,1}=$ | $F$ | $E$ | 1 | $F$ | $F$ | $E$ | 1 | $F$ | $F$ | $E$ | 0 | $E$ | $F$ | $E$ | 0 | $E$ |
| and $K_{2,2}=$ | 1 | $F$ | $F$ | $E$ | 1 | $F$ | $F$ | $E$ | 0 | $E$ | $F$ | $E$ | 0 | $E$ | $F$ | $E$ |
| $K_{3,1}=$ | $E$ | 0 | 1 | $F$ | $E$ | 0 | 1 | $F$ | $F$ | 1 | 0 | $E$ | $F$ | 1 | 0 | $E$ |
| and $K_{3,2}=$ | 1 | $F$ | $E$ | 0 | 1 | $F$ | $E$ | 0 | 0 | $E$ | $F$ | 1 | 0 | $E$ | $F$ | 1 |

## Question 4:

We have seen that DES linear cryptanalysis exploits the sometimes biased input-output relationship of a given substitution box. Give the input-output relationship of substitution boxes $S_{4}$ and $S_{7}$. Which one is better against linear cryptanalysis? How do $S_{4}$ and $S_{7}$ compare to $S_{5}$ (section 3.6.2 of the course notes). Justify your answers (e.g. tables).

